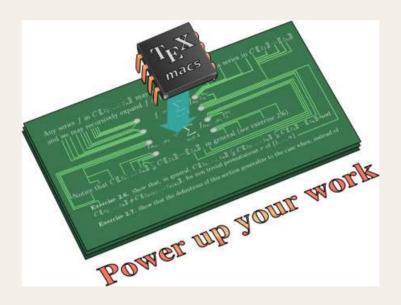
Hardy field solutions to algebraic differential equations

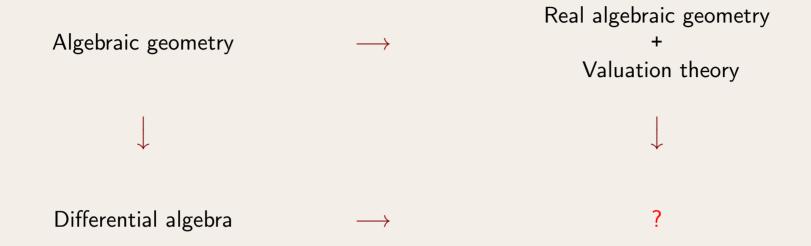


Joris van der Hoeven, Pisa 2007 http://www.TEXmacs.org



A missing subject?





- Hardy fields: Hardy, Rosenlicht, Boshernitzan, Singer, etc.
- Pfaff systems: Khovanskii, Wilkie, van den Dries, Rolin, etc.



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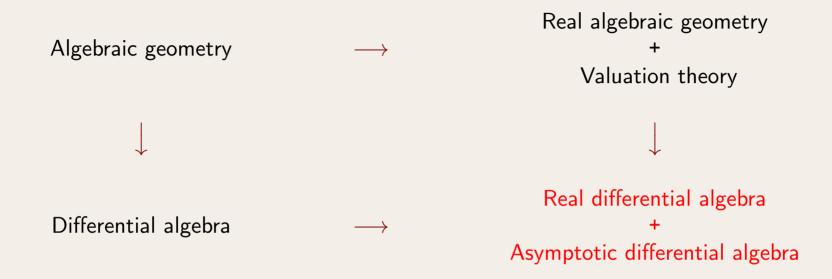


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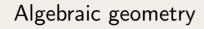




- LNM 1888: Transseries and Real Differential Algebra
- Other work on http://www.math.u-psud.fr/~vdhoeven







Real algebraic geometry

+

Valuation theory

Differential algebra —

Real differential algebra

+

Asymptotic differential algebra



Differential algebra

Sufficiently closed models

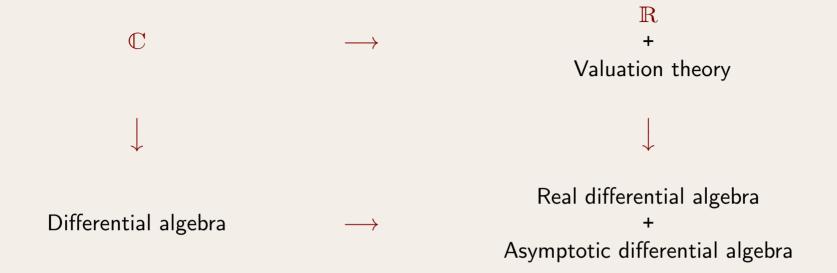




Real differential algebra + Asymptotic differential algebra

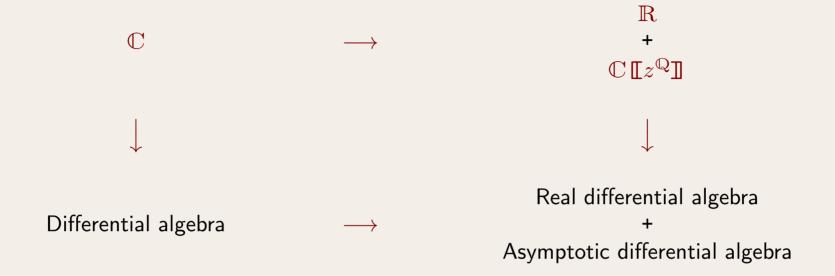






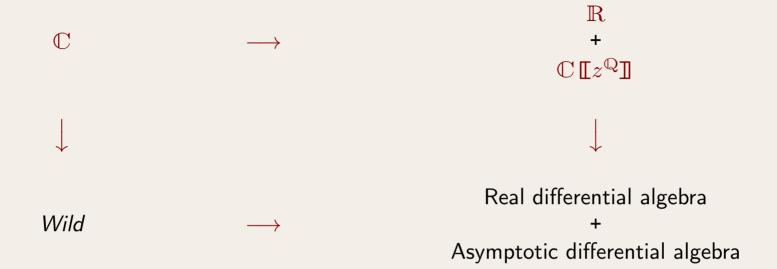






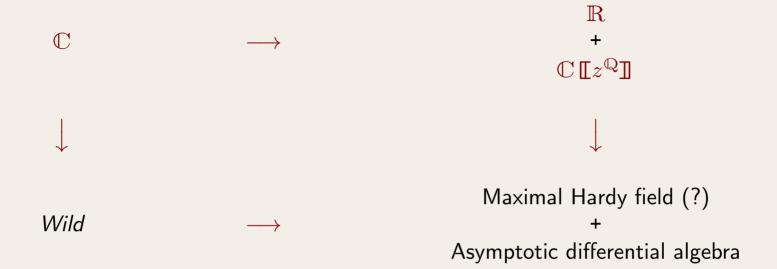






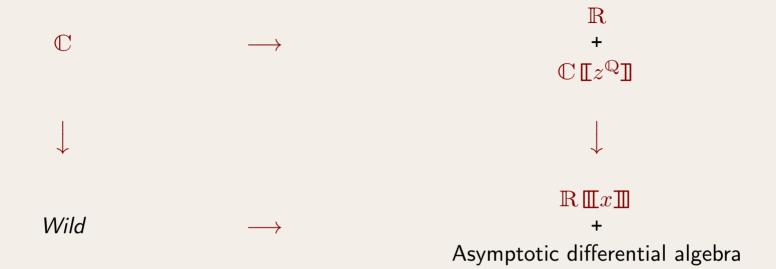
















 \mathbb{C}

 \longrightarrow

 $\mathbb{R} \\ + \\ \mathbb{C} \llbracket z^{\mathbb{Q}} \rrbracket$

Wild

 \longrightarrow

 $\mathbb{R} \, [\![x]\!] \\ + \\ \mathbb{C} \, [\![z]\!]$



What is a transseries?



 $(x \succ 1)$

$$e^{e^x + \frac{e^x}{x} + \frac{e^x}{x^2} + \dots} + \frac{2}{\log x} e^{e^x + \frac{e^x}{x} + \frac{e^x}{x^2} + \dots} + e^{\sqrt{x} + e^{\sqrt{\log x} + e^{\sqrt{\log \log x} + \dots}}} + \dots$$



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- Dahn & Göring
- Écalle



Examples of transseries



$$\frac{1}{1-x^{-1}-x^{-e}} = 1+x^{-1}+x^{-2}+x^{-e}+x^{-3}+x^{-e-1}+\cdots$$

$$\frac{1}{1-x^{-1}-e^{-x}} = 1+\frac{1}{x}+\frac{1}{x^2}+\cdots+e^{-x}+2\frac{e^{-x}}{x}+\cdots+e^{-2x}+\cdots$$

$$-e^x\int\frac{e^{-x}}{x} = \frac{1}{x}-\frac{1}{x^2}+\frac{2}{x^3}-\frac{6}{x^4}+\frac{24}{x^5}-\frac{120}{x^6}+\cdots$$

$$\Gamma(x) = \frac{\sqrt{2\pi}\,e^{x(\log x-1)}}{x^{1/2}}+\frac{\sqrt{2\pi}\,e^{x(\log x-1)}}{12\,x^{3/2}}+\frac{\sqrt{2\pi}\,e^{x(\log x-1)}}{288\,x^{5/2}}+\cdots$$

$$\zeta(x) = 1+2^{-x}+3^{-x}+4^{-x}+\cdots$$

$$\varphi(x) = \frac{1}{x}+\varphi(x^\pi)=\frac{1}{x}+\frac{1}{x^\pi}+\frac{1}{x^{\pi^2}}+\frac{1}{x^{\pi^3}}+\cdots$$

$$\psi(x) = \frac{1}{x}+\psi(e^{\log^2 x})=\frac{1}{x}+\frac{1}{e^{\log^2 x}}+\frac{1}{e^{\log^4 x}}+\frac{1}{e^{\log^8 x}}+\cdots$$



The field ${\mathbb T}$ of grid-based transseries



- $\mathbb{T} = \mathbb{R} \llbracket \mathfrak{T} \rrbracket$, where \mathfrak{T} is a totally ordered monomial group.
- $\mathbb{R} \, \llbracket \mathfrak{T} \rrbracket$: series $f = \sum_{\mathfrak{m} \in \mathfrak{T}} \, f_{\mathfrak{m}} \, \mathfrak{m} \in \mathbb{R} \, \llbracket \mathfrak{T} \rrbracket$ with **grid-based support**:

supp
$$f \subseteq {\{\mathfrak{m}_1, ..., \mathfrak{m}_m\}^* \mathfrak{n}}, \qquad \mathfrak{m}_1, ..., \mathfrak{m}_m \prec 1$$

- T is a totally ordered, real closed field.
- \mathbb{T} is stable under exp, \log , ∂ , \int , \circ and inv.





Theorem. (2000) Given $P \in \mathbb{T}\{F\}$ and $f < g \in \mathbb{T}$ with P(f) P(g) < 0. Then there exists an $h \in \mathbb{T}$ with f < h < g and P(h) = 0.

- 1. Calculus with **cuts** $\hat{f} \in \hat{\mathbb{T}}$.
- 2. Classification of cuts and behaviour of P(f) near a cut.
- 3. Newton polygon method for shrinking interval on which a sign change occurs and whose end-points are cuts.





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Corollary. Any $P \in \mathbb{T}\{F\}$ of odd degree admits a root in \mathbb{T} .





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Example. The following equation admits a solution in \mathbb{T} :

$$\frac{1}{x}f'''(f')^2 f^{24} + e^x (f'')^{27} - \Gamma(\Gamma(\log x)) f^2 = \frac{e^{e^x + x^2}}{\Gamma(e^x + x)}.$$





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Corollary. Any monic $L \in \mathbb{T}[\partial]$ admits a factorization with factors

$$\partial - a$$
 or

$$\partial^2 - (2 \, a + b^{\dagger}) \, \partial + (a^2 + b^2 - a' + a \, b^{\dagger}) = (\partial - (a - b \, \mathbf{i} + b^{\dagger})) \, (\partial - (a + b \, \mathbf{i}))$$



Complex transseries



Theorem. (2001) Every asymptotic differential equation over \mathbb{T} of Newton degree d admits at least d solutions (when counting with multiplicities). Moreover, it suffices to add iterated logarithms to the asymptotic scale.



Complex transseries



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Warning. T is not differentially algebraically closed

$$f^3 + (f')^2 + f = 0$$
$$f^3 + f \neq 0$$

→ Desingularization of vector fields (Cano, Panazzolo, ...)



Complex transseries



Theorem. (2001) Every asymptotic differential equation over \mathbb{T} of Newton degree d admits at least d solutions (when counting with multiplicities). Moreover, it suffices to add iterated logarithms to the asymptotic scale.

Corollary. T is Picard-Vessiot closed.

Remark. ∃ algorithm for computing the solutions of a given equation.

Remark. Zero-test algorithm for polynomials in power series solutions to algebraic differential equations.

Real transseries solutions \rightarrow analytic germs



1: Accelero-summation

2: Transserial Hardy fields

$$\mathbb{T} \ \supseteq \ \mathcal{T} \overset{
ho}{\hookrightarrow} \ \mathcal{G}$$

• \mathcal{G} : ring of infinitely differentiable real germs at $+\infty$.



Real transseries solutions \rightarrow analytic germs



1: Accelero-summation

Advantages	Disadvantages
Canonical after choosing average Preserves composition Classification local vector fields Differential Galois theory	Requires many different tools Not yet written down

2: Transserial Hardy fields

Advantages	Disadvantages
Less hypotheses on coefficients	Not canonical
Might generalize to other models Written down	No preservation of composition





A **transserial Hardy** field is a differential subfield \mathcal{T} of \mathbb{T} , together with a monomorphism $\rho: \mathcal{T} \to \mathcal{G}$ of ordered differential \mathbb{R} -algebras, such that

TH1.
$$\forall f \in \mathcal{T}$$
: supp $f \subseteq \mathcal{T}$.

TH2.
$$\forall f \in \mathcal{T}$$
: $f_{\prec} \in \mathcal{T}$.

$$f_{\prec} = \sum_{\mathfrak{m} \prec 1} f_{\mathfrak{m}} \mathfrak{m}$$

TH3.
$$\exists d \in \mathbb{Z}$$
: $\forall \mathfrak{m} \in \mathfrak{T} \cap \mathcal{T}$: $\log \mathfrak{m} \in \mathcal{T} + \mathbb{R} \log_d x$.

TH4.
$$\mathfrak{T} \cap \mathcal{T}$$
 is stable under taking real powers.

TH5.
$$\forall f \in \mathcal{T}^{>}$$
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Elementary extensions



Definitions. \mathcal{T} transserial Hardy field, $f \in \mathbb{T}$, $\hat{f} \in \mathcal{G}$

$$f \sim \hat{f} \iff (\exists \varphi \in \mathcal{T}: \ f \sim_{\mathbb{T}} \varphi \sim_{\mathcal{G}} \hat{f})$$

$$f \ asympt. \ equiv. \ \text{to} \ \hat{f} \ \text{over} \ \mathcal{T} \iff (\forall \varphi \in \mathcal{T}: \ f - \varphi \sim \hat{f} - \varphi)$$

$$f \ diff. \ equiv. \ \text{to} \ \hat{f} \ \text{over} \ \mathcal{T} \iff (\forall P \in \mathcal{T} \{F\}: \ P(f) = 0 \Leftrightarrow P(\hat{f}) = 0)$$

Lemma. Let $f \in \mathbb{T} \setminus \mathcal{T}$ and $\hat{f} \in \mathcal{G} \setminus \mathcal{T}$ be such that

- i. f is a serial cut over \mathcal{T} .
- ii. f and \hat{f} are asymptotically equivalent over \mathcal{T} .
- iii. f and \hat{f} are differentially equivalent over \mathcal{T} .
- Then $\exists !$ transserial Hardy field extension $\rho : \mathcal{T}\langle f \rangle \to \mathcal{G}$ with $\rho(f) = \hat{f}$.



Basic extension theorems



Theorem. Let \mathcal{T} be a transserial Hardy field. Then its real closure \mathcal{T}^{rcl} admits a unique transserial Hardy field structure which extends the one of \mathcal{T} .

Theorem. Let \mathcal{T} be a transserial Hardy field and let $\varphi \in \mathcal{T}_{\succ}$ be such that $e^{\varphi} \notin \mathcal{T}$. Then the set $\mathcal{T}(e^{\mathbb{R}\varphi})$ carries the structure of a transserial Hardy field for the unique differential morphism $\rho: \mathcal{T}(e^{\mathbb{R}\varphi}) \to \mathcal{G}$ over \mathcal{T} with $\rho(e^{\lambda \varphi}) = e^{\lambda \rho(\varphi)}$ for all $\lambda \in \mathbb{R}$.

Theorem. Let \mathcal{T} be a transserial Hardy field of depth $d < \infty$. Then $\mathcal{T}((\log_d x)^\mathbb{R})$ carries the structure of a transserial Hardy field for the unique differential morphism $\rho: \mathcal{T}((\log_d x)^\mathbb{R}) \to \mathcal{G}$ over \mathcal{T} with $\rho((\log_d x)^\lambda) = (\log_d x)^\lambda$ for all $\lambda \in \mathbb{R}$.



Differential equations (main ideas)



Step 1. A given algebraic differential equation

$$f^2 - f' + \frac{x}{e^x} = 0$$

Step 2. Put equation in integral form

$$f = \int \left(\frac{x}{e^x} + f^2\right)$$

Step 3. Integral transseries solution



Differential equations (main ideas)



Step 1. A given algebraic differential equation

$$f^2 - e^x f' + \frac{e^{2x}}{x} = 0$$

Step 2. Put equation in integral form

$$f = \int \left(\frac{e^x}{x} + \frac{f^2}{e^x} \right)$$

Step 3. Integrate from a fixed point $x_0 < \infty$



Differential equations (main ideas)



Step 1. A general algebraic differential equation

$$P(f) = 0$$

Step 2. Equation in split-normal form

$$(\partial - \varphi_1) \cdots (\partial - \varphi_r) f = P(f)$$
 with $P(f)$ small

Attention: $\varphi_1, ..., \varphi_r \in \mathcal{T}[i]$, even though $(\partial - \varphi_1) \cdots (\partial - \varphi_r) \in \mathcal{T}[\partial]$.

Step 3. Solve the split-normal equation using the fixed-point technique.



Continuous right-inverses (first order)



Lemma. The operator $J = (\partial - \varphi)_{x_0}^{-1}$, defined by

$$(Jf)(x) = \begin{cases} e^{\Phi(x)} \int_{-\infty}^{x} e^{-\Phi(t)} f(t) dt & (repulsive \ case) \\ e^{\Phi(x)} \int_{-x_0}^{x} e^{-\Phi(t)} f(t) dt & (attractive \ case) \end{cases}$$

and

$$\Phi(x) = \begin{cases} \int_{-\infty}^{x} \varphi(t) dt & (repulsive \ case) \\ \int_{-x_0}^{x} \varphi(t) dt & (attractive \ case) \end{cases}$$

is a continuous right-inverse of $L = \partial - \varphi$ on $\mathcal{G}^{\bowtie}[i]$, with

$$|||J||_{x_0} \leqslant \left\| \frac{1}{\operatorname{Re} \varphi} \right\|_{x_0}$$



Continuous right-inverses (higher order)



Lemma. Given a split-normal operator

$$L = (\partial - \varphi_1) \cdots (\partial - \varphi_r), \tag{1}$$

with a factorwise right-inverse $L^{-1} = J_r \cdots J_1$, the operator

$$\mathfrak{v}^{\nu}J_{r}\cdots J_{1}:\mathcal{G}_{x_{0}}^{\preccurlyeq}[\mathrm{i}] \rightarrow \mathcal{G}_{x_{0};r}^{\preccurlyeq}[\mathrm{i}]$$

is a continuous operator for every $\nu > r \sigma_L$. Here $\mathcal{G}_{x_0;r}^{\preccurlyeq}[i]$ carries the norm

$$||f||_{x_0;r} = \max\{||f||_{x_0},...,||f^{(r)}||_{x_0}\}.$$

Lemma. If $L \in \mathcal{T}[\partial]$ and the splitting (1) (formally) preserves realness, then $J_r \cdots J_1$ preserves realness in the sense that it maps $\mathcal{G}_{x_0}^{\prec}$ into itself.



Non-linear equations



Theorem. Consider a split-monic equation

$$Lf = P(f), \quad f \prec 1,$$

and let ν be such that r $\sigma_L < \nu < v_P$. Then for any sufficiently large x_0 , there exists a continuous factorwise right-inverse $J_{r, \ltimes v^{\nu}} \cdots J_{1, \ltimes v^{\nu}}$ of $L_{\ltimes v^{\nu}}$, such that the operator

$$\Xi: f \longmapsto (J_r \cdots J_1)(P(f))$$

admits a unique fixed point

$$f = \lim_{n \to \infty} \Xi^{(n)}(0) \in \mathcal{B}(\mathcal{G}_{x_0;r}^{\preccurlyeq}, \frac{1}{2}).$$



Preservation of asymptotics



Theorem. Let \mathcal{T} be a transserial Hardy field of span $\mathfrak{v} \succeq e^x$. Consider a monic split-normal quasi-linear equation

$$Lf = P(f), \quad f < 1, \tag{2}$$

over \mathcal{T} without solutions in \mathcal{T} . Assume that one of the following holds:

- a) \mathcal{T} is (1,1,1)-differentially closed in $\mathbb{T}_{\underline{\prec} v}$ and (2) is first order.
 - i.e. \mathcal{T} is closed under the resolution of linear first order equations.
- b) $\mathcal{T}[i]$ is (1,1,1)-differentially closed in $\mathbb{T}[i]_{\ll_{\mathfrak{v}}}$.

Then there exist solutions $f \in \mathcal{G}$ and $\tilde{f} \in \hat{\mathcal{T}}$ to (2), such that f and \tilde{f} are asymptotically equivalent over \mathcal{T} .

(A)

First order extensions



Lemma. Let $L = \partial - \varphi \in \mathcal{T}[\partial]$ be a normal operator. Let $\tilde{f} \in \hat{\mathcal{T}}^{\preccurlyeq}$ and $g \in \mathcal{T}^{\preccurlyeq}$ be such that \tilde{f} is transcendental over \mathcal{T} and $L \tilde{f} = g$. Then there exists an $f \in \mathcal{G}^{\preccurlyeq}$ with Lf = g, such that f and \tilde{f} are both differentially and asymptotically equivalent over \mathcal{T} .

Theorem. Let \mathcal{T} be a transserial Hardy field. Let $\mathcal{T}^{\mathrm{fo}} \supseteq \mathcal{T}$ be the smallest differential subfield of \mathbb{T} , such that for any $P \in \mathcal{T}^{\mathrm{fo}} \{F\}^{\neq}$ with $r_P \leqslant 1$ and $f \in \mathbb{T}$ we have $P(f) = 0 \Rightarrow f \in \mathcal{T}^{\mathrm{fo}}$. Then the transserial Hardy field structure of \mathcal{T} can be extended to $\mathcal{T}^{\mathrm{fo}}$.

Proof. As long as $\mathcal{T}^{\text{fo}} \neq \mathcal{T}$:

- Close off under exp, log and algebraic equations.
- Choose $P \in \mathcal{T}\{F\}^{\neq}$, $r_P = 1$, $f \in \mathbb{T}$, P(f) = 0 such that P has minimal "complexity" (r_P, d_P, t_P) and apply the lemma.



Higher order extensions



Lemma. Let $L = \partial - \varphi \in \mathcal{T}[i][\partial]$ be a normal operator. Let $\tilde{f} \in \hat{\mathcal{T}}[i]^{\preccurlyeq}$ and $g \in \mathcal{T}[i]^{\preccurlyeq}$ be such that $\operatorname{Re} \tilde{f}$ has order 2 over \mathcal{T} and $L \tilde{f} = g$. Then there exists an $f \in \mathcal{G}^{\preccurlyeq}[i]$ with Lf = g, such that $\operatorname{Re} f$ and $\operatorname{Re} \tilde{f}$ are both differentially and asymptotically equivalent over \mathcal{T} .

Theorem. Let \mathcal{T} be a transserial Hardy field. Let $\mathcal{T}^{\mathrm{dalg}} \supseteq \mathcal{T}$ be the smallest differential subfield of \mathbb{T} , such that for any $P \in \mathcal{T}^{\mathrm{dalg}}\{F\}^{\neq}$ and $f \in \mathbb{T}$ we have $P(f) = 0 \Rightarrow f \in \mathcal{T}^{\mathrm{dalg}}$. Then the transserial Hardy field structure of \mathcal{T} can be extended to $\mathcal{T}^{\mathrm{dalg}}$.



Applications



Corollary. There exists a transserial Hardy field \mathcal{T} , such that for any $P \in \mathcal{T}\{F\}$ and $f, g \in \mathcal{T}$ with f < g and P(f) P(g) < 0, there exists a $h \in \mathcal{T}$ with f < h < g and P(h) = 0.

Corollary. There exists a transserial Hardy field \mathcal{T} , such that $\mathcal{T}[i]$ is weakly differentially closed.

Corollary. There exists a differentially Henselian transserial Hardy field \mathcal{T} , i.e., such that any quasi-linear differential equation over \mathcal{T} admits a solution in \mathcal{T} .



A partial inverse



Theorem. Let \mathcal{T} be a transserial Hardy field and \mathcal{H} a differentially algebraic Hardy field extension of \mathcal{T} , such that \mathcal{H} is differentially Henselian and stable under exponentiation. Then there exists a transserial Hardy field structure on \mathcal{H} which extends the structure on \mathcal{T} .

Corollary. Let \mathcal{T} be a transserial Hardy field and \mathcal{H} a differentially algebraic Hardy field extension of \mathcal{T} , such that \mathcal{H} is differentially Henselian. Assume that \mathcal{H} admits no non-trivial algebraically differential Hardy field extensions. Then \mathcal{H} satisfies the differential intermediate value property.

Theorem. (Boshernitzan 1987) Any solution of the equation

$$f'' + f = e^{x^2}$$

is contained in a Hardy field. However, none of these solutions is contained in the intersection of all maximal Hardy fields.

Open questions



- 1. Embeddability of Hardy fields in differentially Henselian Hardy fields.
- 2. Do maximal Hardy fields satisfy the intermediate value property?
- 3. Restricted analytic (instead of algebraic) differential equations.
- 4. Preservation of composition:
 - a. $f(x+\varepsilon)$, small ε : expand.
 - b. $f(qx+\varepsilon)$: expand, but more intricate.
 - c. $f(\varphi(x)), \varphi \succ x$: abstract nonsense.