

Lesson 4 — Newton polygon method

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Definition

A **support type** for a monomial monoid \mathfrak{M} is a subset $\mathcal{S}(\mathfrak{M}) \subseteq \mathcal{P}(\mathfrak{M})$ such that

T1. Every $\mathcal{S} \in \mathcal{S}(\mathfrak{M})$ is well-based.

T2. If $m \in \mathfrak{M}$, then $\{m\} \in \mathcal{S}(\mathfrak{M})$.

T3. If $\mathcal{S} \in \mathcal{S}(\mathfrak{M})$ and $\mathcal{T} \subseteq \mathcal{S}$, then $\mathcal{T} \in \mathcal{S}(\mathfrak{M})$.

T4. If $\mathcal{S}, \mathcal{T} \in \mathcal{S}(\mathfrak{M})$, then $\mathcal{S} \cup \mathcal{T} \in \mathcal{S}(\mathfrak{M})$.

T5. If $\mathcal{S}, \mathcal{T} \in \mathcal{S}(\mathfrak{M})$, then $\mathcal{S} \mathcal{T} := \{mn : m \in \mathcal{S}, n \in \mathcal{T}\} \in \mathcal{S}(\mathfrak{M})$.

T6. If $\mathcal{S} \in \mathcal{S}(\mathfrak{M})$ and $\mathcal{S} < 1$, then $\mathcal{S}^* := \{m_1 \cdots m_n : m_1, \dots, m_n \in \mathcal{S}\} \in \mathcal{S}(\mathfrak{M})$.

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Let \mathcal{S} be a map that associates a support type $\mathcal{S}(\mathfrak{M})$ for \mathfrak{M} to any monomial monoid \mathfrak{M} .

We say that \mathcal{S} is a **support type** if:

ST. For every strictly increasing morphism $\varphi: \mathfrak{M} \rightarrow \mathfrak{N}$ and $\mathcal{S} \in \mathcal{S}(\mathfrak{M})$, we have $\varphi(\mathcal{S}) \in \mathcal{S}(\mathfrak{N})$.

\mathcal{P} -based series

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We denote by $R[[\mathfrak{M}]]_{\mathcal{P}}$ the set of all such series.

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A well-based family $(f_i)_{i \in \mathfrak{M}} \in R[[\mathfrak{M}]]_{\mathcal{P}}$ is **\mathcal{P} -based** if $\bigcup_{i \in I} \text{supp } f_i \in \mathcal{P}(\mathfrak{M})$.

Then $\sum_{i \in I} f_i \in R[[\mathfrak{M}]]_{\mathcal{P}}$. This defines “the natural” strong summation on $R[[\mathfrak{M}]]_{\mathcal{P}}$.

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Proposition

a) $R[[\mathfrak{M}]]_{\mathcal{P}}$ is a ring.

b) If R is a field and \mathfrak{M} a totally ordered group, then $R[[\mathfrak{M}]]_{\mathcal{P}}$ is a field.

Well-based supports.

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$$\zeta(x) = 1 + 2^{-x} + 3^{-x} + \dots = 1 + e^{-(\log 2)x} + e^{-(\log 3)x} + \dots \notin \mathbb{R}[[e^{-\mathbb{R}x}]]_{\mathcal{S}}.$$

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Intersections. Let \mathcal{S} and \mathcal{T} be support types.

$$(\mathcal{S} \cap \mathcal{T})(\mathfrak{M}) = \mathcal{S}(\mathfrak{M}) \cap \mathcal{T}(\mathfrak{M}).$$

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We say that $\mathcal{S} \subseteq \mathfrak{M}$ is **grid-based** if there exist finite sets $\mathfrak{F} \subseteq \mathfrak{M}$ and $\mathcal{E} \subseteq \mathfrak{M}^{<1}$ with

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Note. If \mathfrak{M} is a totally ordered group, then \mathcal{F} can be taken to be a singleton.

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If $\mathcal{S} \subseteq \mathfrak{M}^{<1}$ is grid-based, then there is a finite $\mathcal{E} \subseteq \mathfrak{M}^{<1}$ with $\mathcal{S} \subseteq \mathcal{E}^*$ (whence $\mathcal{S}^* \subseteq \mathcal{E}^*$).

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Proof. Let $\mathfrak{F} \subseteq \mathfrak{M}$ and $\mathcal{G} \subseteq \mathfrak{M}^{<1}$ be finite with $\mathcal{S} \subseteq \mathfrak{F} \mathcal{G}^*$.

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Let $\mathfrak{H}_f \subseteq \mathfrak{M}^{<1}$ be a finite set of generators. Note that $(f\mathcal{G}^*) \cap \mathfrak{M}^{<1} \subseteq \mathfrak{H}_f\mathcal{G}^*$.

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Now it suffices to take $\mathcal{E} := \mathcal{G} \cup \bigcup_{f \in \mathfrak{F}} \mathfrak{H}_f$.

□

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\mathcal{G} -based families are called **grid-based families**. Etc.

Proposition

For any $f \in R[[\mathfrak{M}]]$, there exist power series $\check{f}_1, \dots, \check{f}_l \in R[[z_1, \dots, z_k]]$, monomials $f_1, \dots, f_l \in \mathfrak{M}$ and $e_1, \dots, e_k \in \mathfrak{M}^{<1}$ with

$$f = \sum_{1 \leq i \leq l} (\check{f}_i \circ (e_1, \dots, e_k)) f_i.$$

Proposition

Assume that \mathfrak{M} is a totally ordered group.

For any $f \in R[[\mathfrak{M}]]$, there exists a Laurent series $\check{f} \in R((z_1, \dots, z_k))$ and $e_1, \dots, e_k \in \mathfrak{M}^{<1}$ with

$$f = \check{f} \circ (e_1, \dots, e_k). \quad (\star)$$

Here $(g z_1^{i_1} \cdots z_k^{i_k}) \circ (e_1, \dots, e_k) := (g \circ (e_1, \dots, e_k)) e_1^{i_1} \cdots e_k^{i_k}$ for any $g \in R[[z_1, \dots, z_k]]$, $i_1, \dots, i_k \in \mathbb{Z}$.

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We call (\star) a **Cartesian representation** of f .

Definition

Let \mathcal{L} be a collection of subsets $\mathcal{L}_k \subseteq R[[z_1, \dots, z_k]]$ for $k \in \mathbb{N}$, such that

L1. $z_i \in \mathcal{L}_k$ for $i = 1, \dots, k$.

L2. \mathcal{L}_k is an R -subalgebra of $R[[z_1, \dots, z_k]]$.

L3. For any $f \in \mathcal{L}_k$ with $z_1 \mid f$, we have $z_1^{-1} f \in \mathcal{L}_k$.

L4. Given $f \in \mathcal{L}_k$ and $g_1, \dots, g_k \in \mathcal{L}_1^{<1}$, we have $f \circ (g_1, \dots, g_k) \in \mathcal{L}_1$.

L5. Given $f \in \mathcal{L}_{k+1}$ with $f(0, \dots, 0) = 0$ and $(\partial f / \partial z_{k+1})(0, \dots, 0) = 1$, the unique $\varphi \in R[[z_1, \dots, z_k]]$ with $f \circ (z_1, \dots, z_k, \varphi) = 0$ is in \mathcal{L}_k .

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- $\mathcal{L}_k = K[[z_1, \dots, z_k]]^{\text{alg}}$, algebraic power series, K any field.

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- $\mathcal{L}_k = K[[z_1, \dots, z_k]]^{\text{alg}}$, algebraic power series, K any field.
- $\mathcal{L}_k = K[[z_1, \dots, z_k]]^{\text{dalg}}$, d-algebraic power series, K any field with $\text{char } K = 0$.

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for some $\check{f} \in \mathcal{L}_k z_1^{\mathbb{Z}} \cdots z_k^{\mathbb{Z}}$ and $\mathbf{e}_1, \dots, \mathbf{e}_k \in \mathfrak{M}^{<1}$.

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Theorem

If K is a field, then so is $K[[\mathfrak{M}]]_{\mathcal{L}}$ is a field. Moreover, if \mathfrak{M} has \mathbb{Q} -powers, then

- If K is algebraically closed and of characteristic zero, then so is $K[[\mathfrak{M}]]_{\mathcal{L}}$.
- If K is real closed, then so is $K[[\mathfrak{M}]]_{\mathcal{L}}$.

K algebraically closed field

Γ divisible totally ordered *abelian* group: $(\forall \gamma \in \Gamma) (\forall n \in \mathbb{N}^{>0}) (\exists \alpha \in \Gamma) n\alpha = \gamma$

z^Γ corresponding monomial group, $z^\alpha \preceq z^\beta \Leftrightarrow \alpha \succeq \beta$.

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Our goal

Given $P \in K[[z^\Gamma]][Y] \setminus K[[z^\Gamma]]$, compute the solutions in $K[[z^\Gamma]]$ of

$$P(y) = 0.$$

K algebraically closed field

Γ divisible totally ordered *abelian* group: $(\forall \gamma \in \Gamma) (\forall n \in \mathbb{N}^{>0}) (\exists \alpha \in \Gamma) n\alpha = \gamma$

z^Γ corresponding monomial group, $z^\alpha \preceq z^\beta \Leftrightarrow \alpha \geq \beta$.

Our goal

Given $P \in K[[z^\Gamma]][Y] \setminus K[[z^\Gamma]]$ and $\gamma \in \Gamma$, compute the solutions in $K[[z^\Gamma]]$ of

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We may replace $K[[z^\Gamma]]$ by $K[[z^\Gamma]]_{\mathcal{F}}$ or $K[[z^\Gamma]]_{\mathcal{G}}$.

$$P_d y^d + \cdots + P_0 = 0, \quad (y < z^\gamma). \quad (\star)$$

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If y satisfies (\star) , it follows that there exists a $j \neq i$ with

$$P_i y^i = P_j y^j \succcurlyeq P_k y^k, \quad \text{for all } k.$$

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$$P_i y^i = P_j y^j \geq P_k y^k, \quad \text{for all } k.$$

Setting $z^{\pi_k} := \partial_{P_k}$ for $k = 0, \dots, d$, and $z^\nu := \partial_y$, this means that there exist $i \neq j$ with

$$\nu > \gamma, \quad \pi_i + i\nu = \pi_j + j\nu \leq \pi_k + k\nu, \quad \text{for all } k.$$

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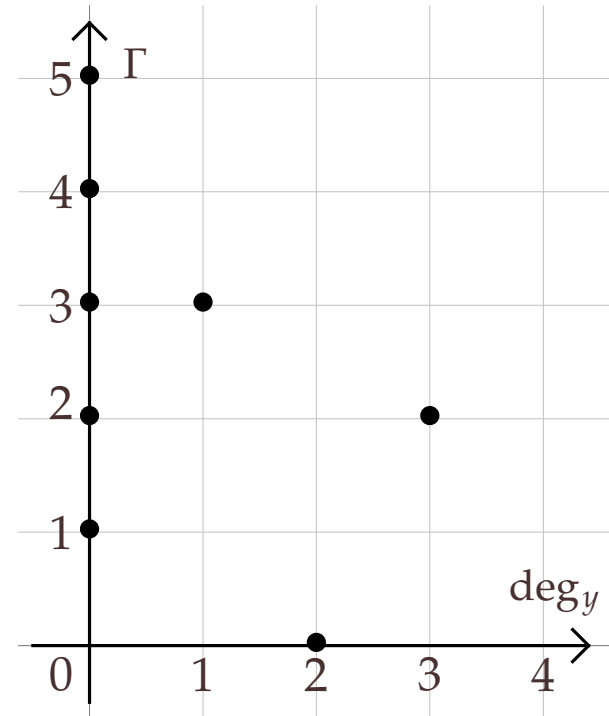
$$\nu > \gamma, \quad \pi_i + i\nu = \pi_j + j\nu \leq \pi_k + k\nu, \quad \text{for all } k.$$

We call z^ν a **starting monomial** for the equation (\star) .

Newton polygons

12/27

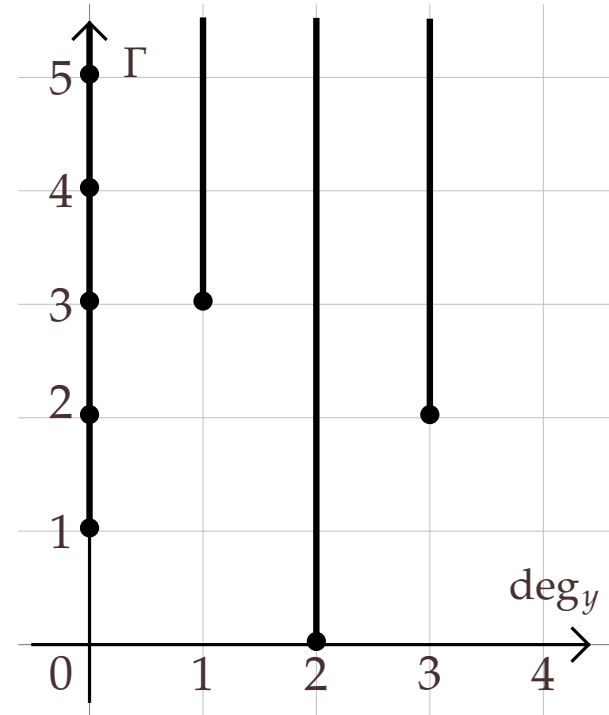
$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$



Newton polygons

12/27

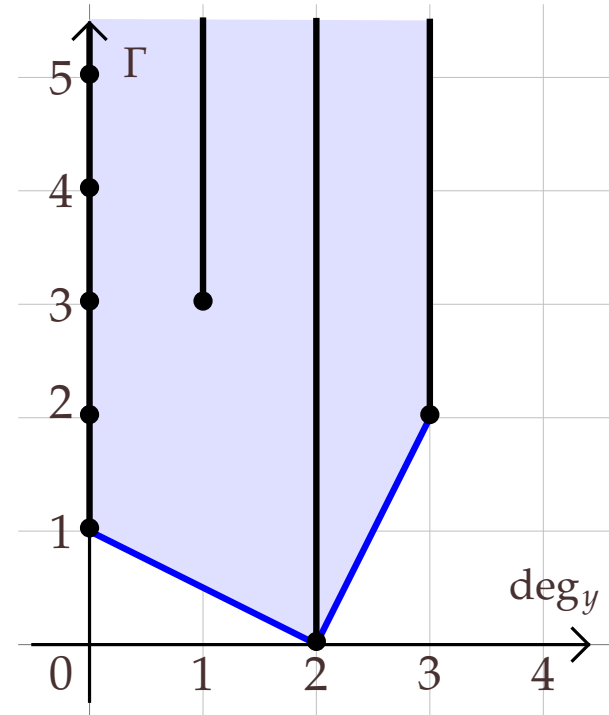
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Newton polygons

12/27

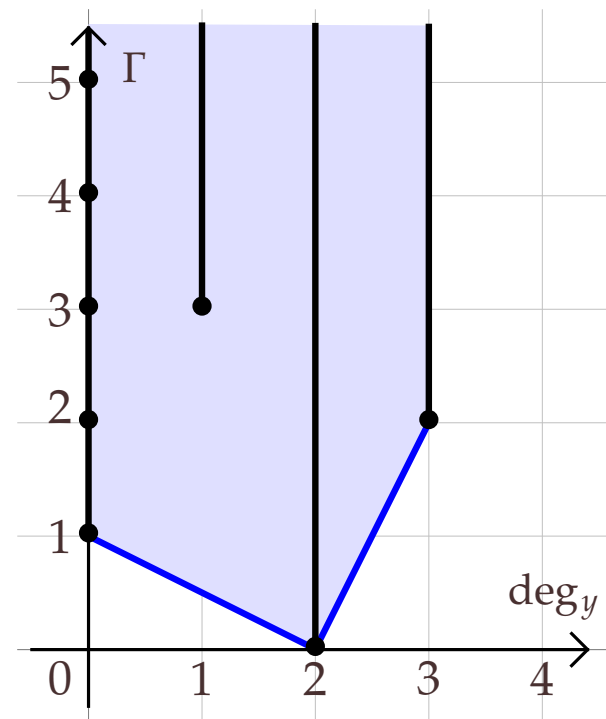
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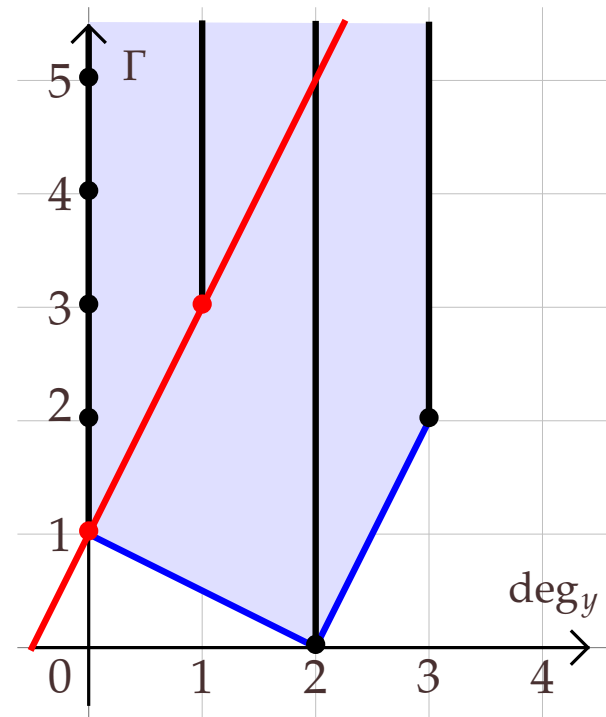


Newton polygons

$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Starting monomials $z^\nu \asymp y$?

- $P_0 \asymp P_1 y \implies z = z^{3+\nu} \implies \nu = -2$

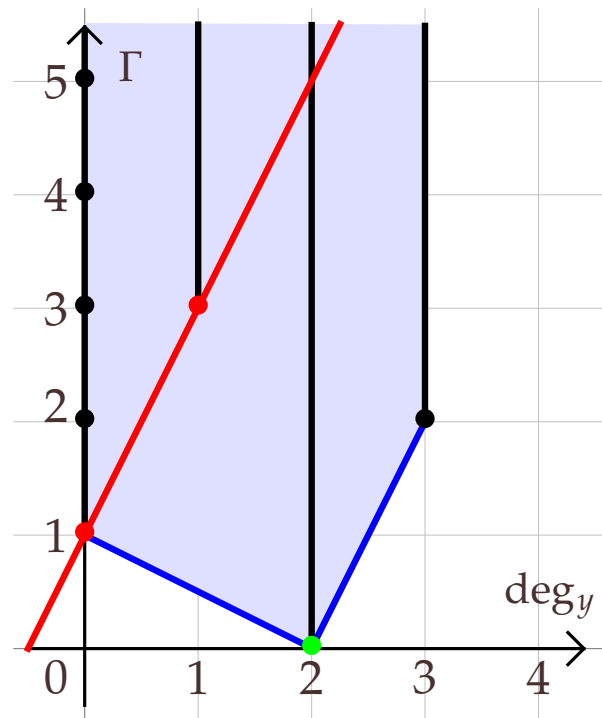


$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Starting monomials $z^v \asymp y$?

- $P_0 \asymp P_1 y \implies z = z^{3+v} \implies v = -2$

But then $P_2 y^2 \asymp z^{0+2v} = z^{-4} \succ z \asymp P_0$

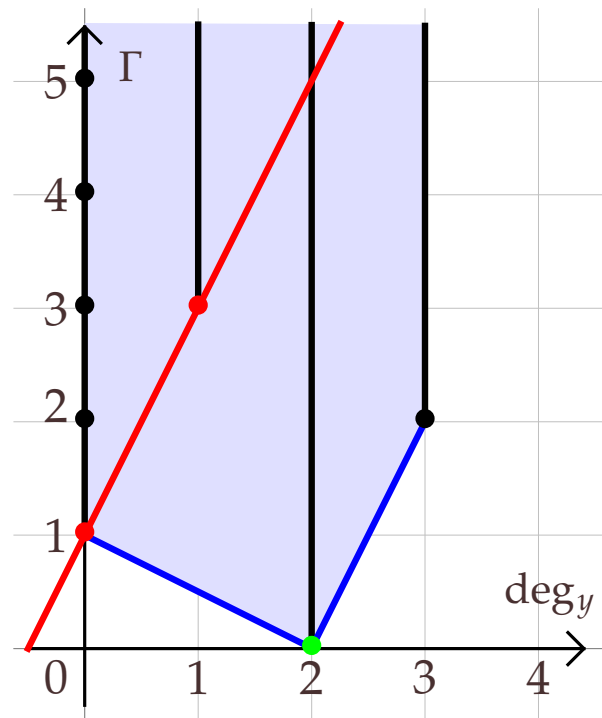


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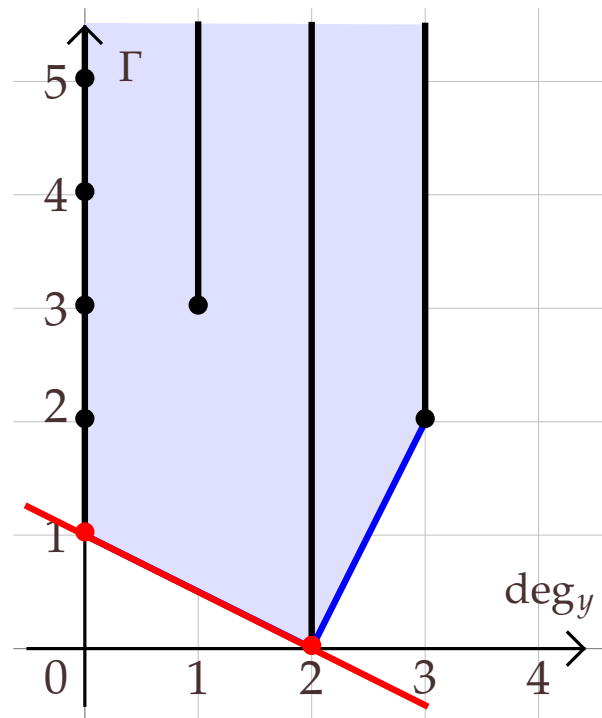
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- $P_0 \asymp P_2 y^2 \implies z = z^{0+2\nu} \implies \nu = 1/2$



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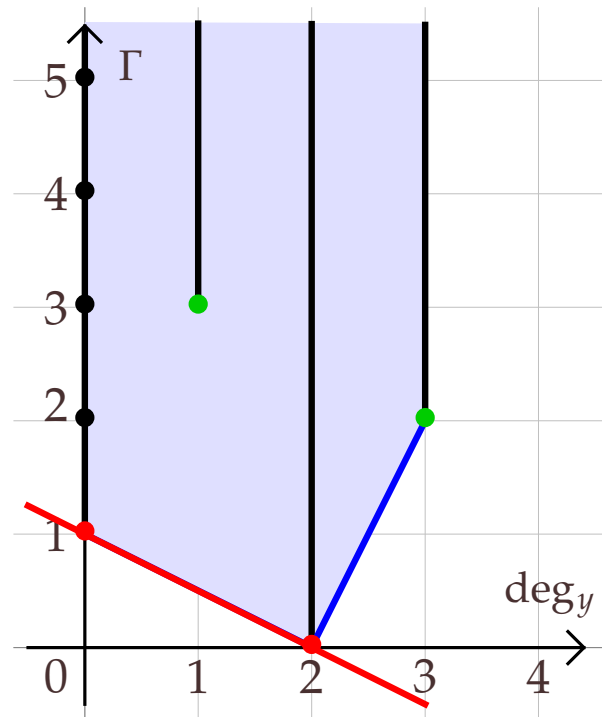
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OK, since $P_1 y \asymp z^{3+\nu} = z^{3^{1/2}} \ll z \asymp P_0$

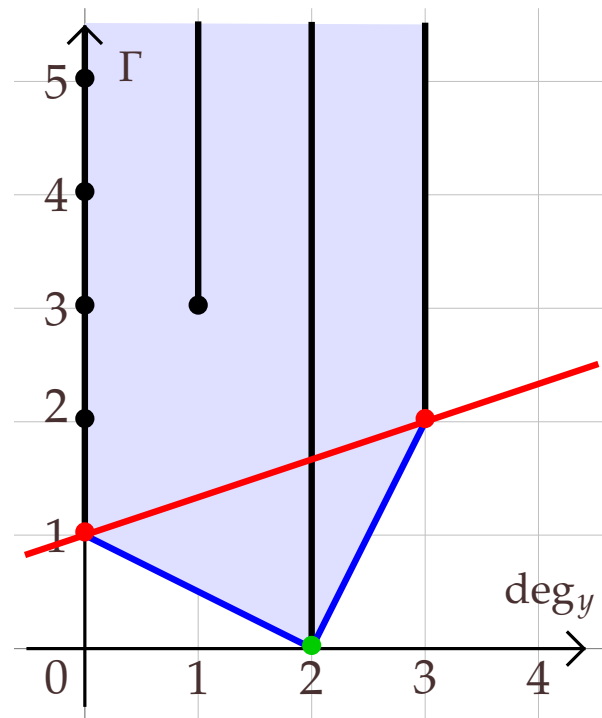
$$P_3 y^3 \asymp z^{2+3\nu} = z^{3^{1/2}} \ll z \asymp P_0$$



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Starting monomials $z^v \asymp y$?

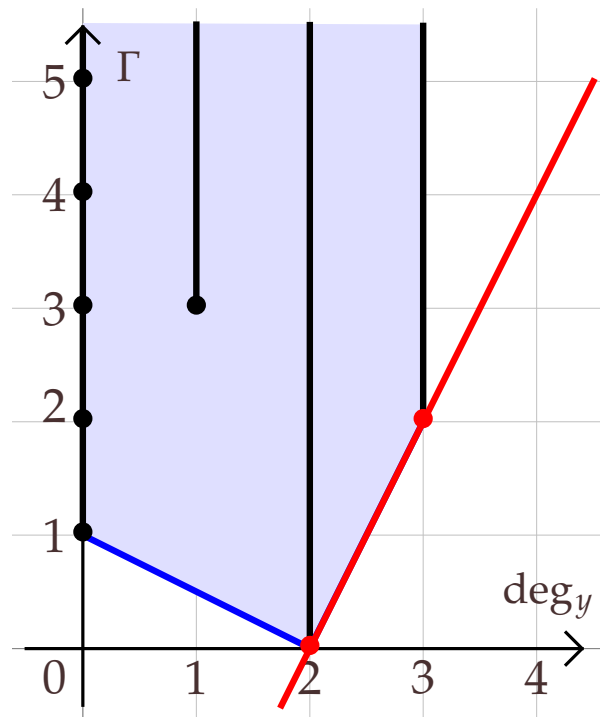
- $P_0 \asymp P_1 y \implies z = z^{3+\nu} \implies \nu = -2$
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- $P_0 \asymp P_2 y^2 \implies z = z^{0+2\nu} \implies \nu = 1/2$
 OK, since $P_1 y, P_3 y^2 \preceq P_0$
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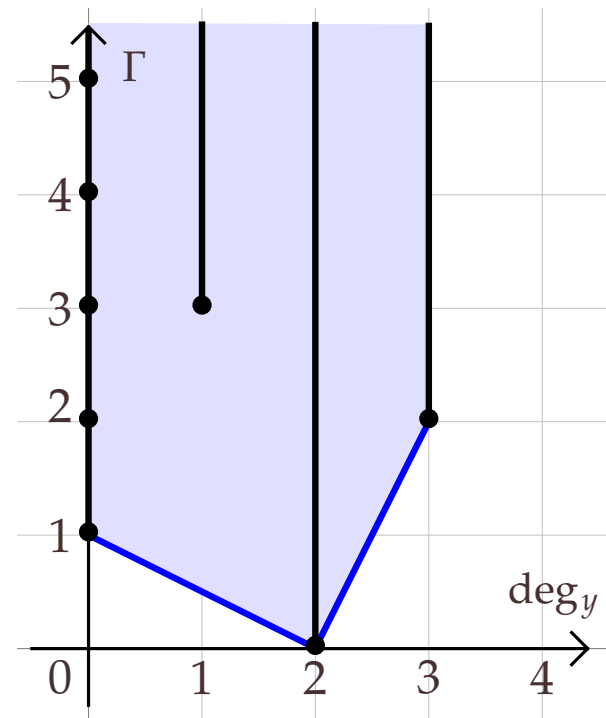


Newton polygons

$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Starting monomials $z^\nu \asymp y$?

- $\nu = 1/2$
- $\nu = -2$



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Consider the starting monomial $z^{1/2}$.

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If $y \sim cz^{1/2}$, then

$$\begin{aligned}5z^2y^3 &< z \\ y^2 &\sim c^2z \\ 3z^2y &< z \\ -\frac{z}{1-z} &\sim -z\end{aligned}$$

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$$c^2 - 1 = 0,$$

whence

$$c = 1 \vee c = -1.$$

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We call $z^{1/2}$ and $-z^{1/2}$ **starting terms** for the equation.

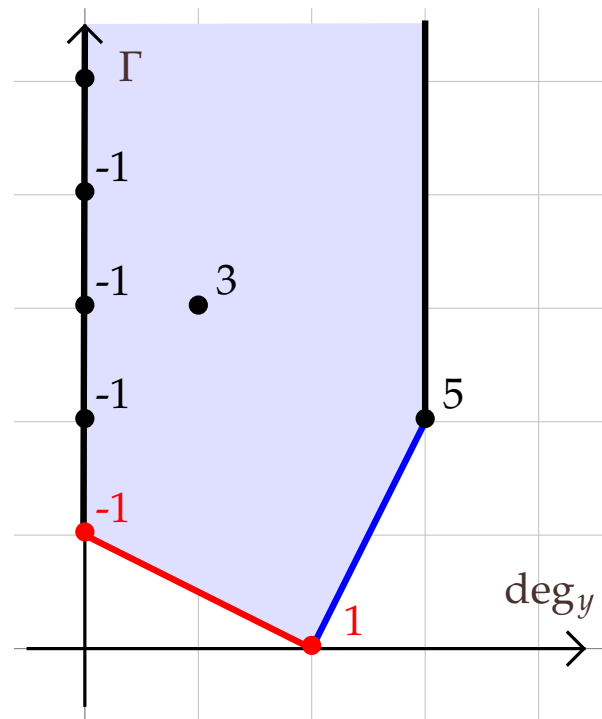
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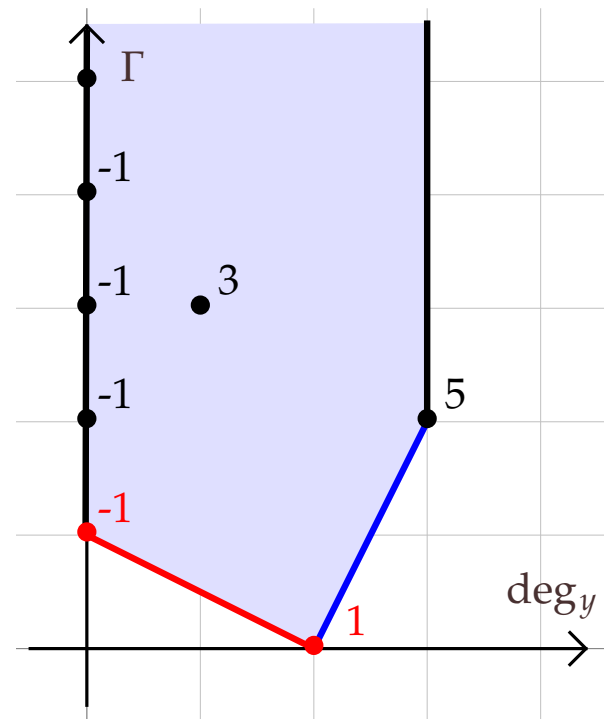
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$N(c) = c^2 - 1$ is the **Newton polynomial** for $z^{1/2}$.



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

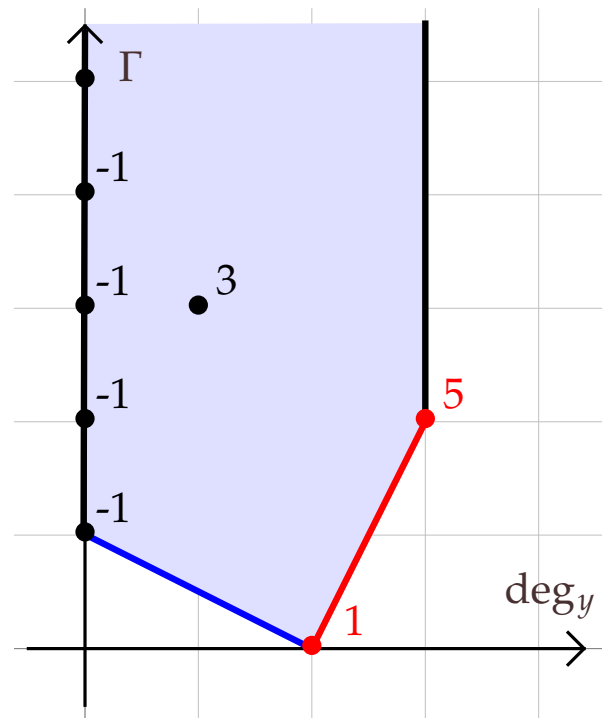
If $y \sim cz^{-2}$, then $c \neq 0$ and

$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = (5c^3 + c^2)z^{-4} + o(z^{-4}).$$

If the right-hand side vanishes, then

$$5c^3 + c^2 = 0.$$

$N(c) = 5c^3 + c^2$ is the **Newton polynomial** for z^{-2} .



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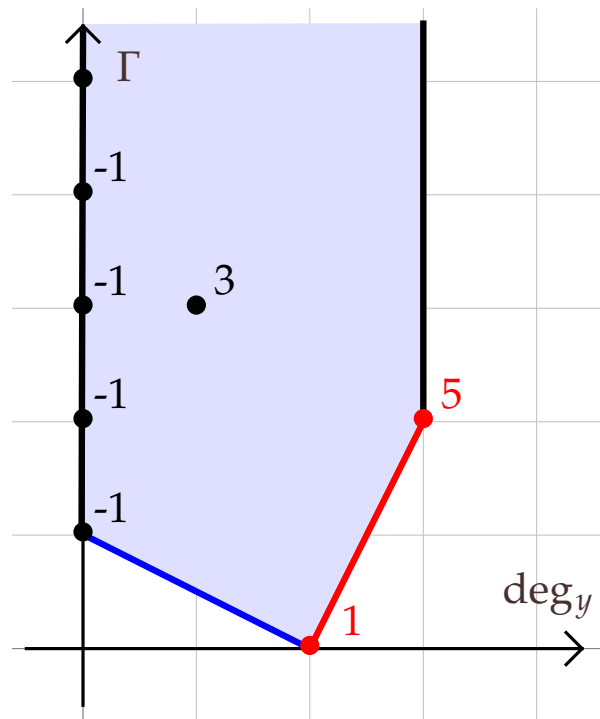
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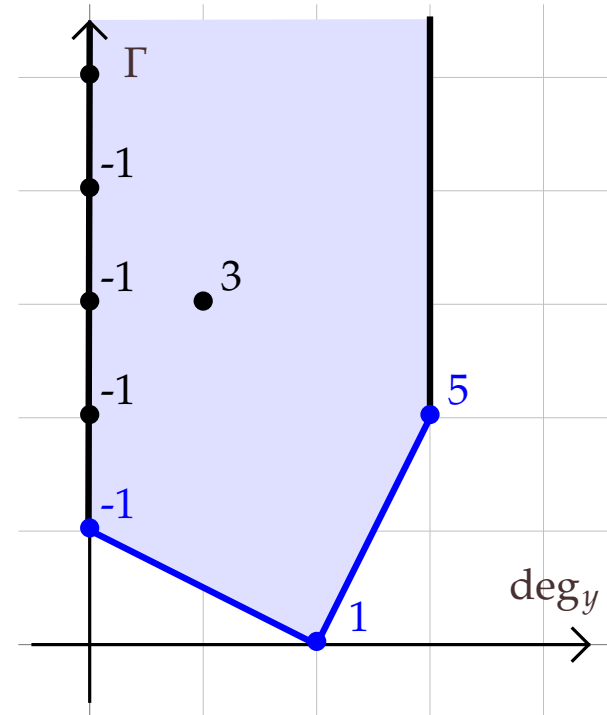
$-1/5z^{-2}$ is a starting term for the equation.



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

The starting terms for the equation are:

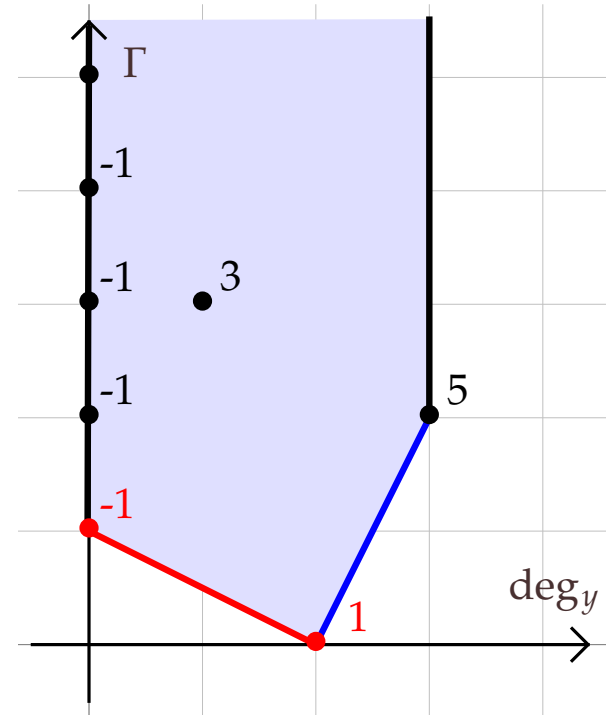
- $z^{1/2}$
- $-z^{1/2}$
- $-1/5z^{-2}$



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Assume $y \sim z^{1/2}$ and perform the change of variables

$$y = z^{1/2} + \tilde{y}$$

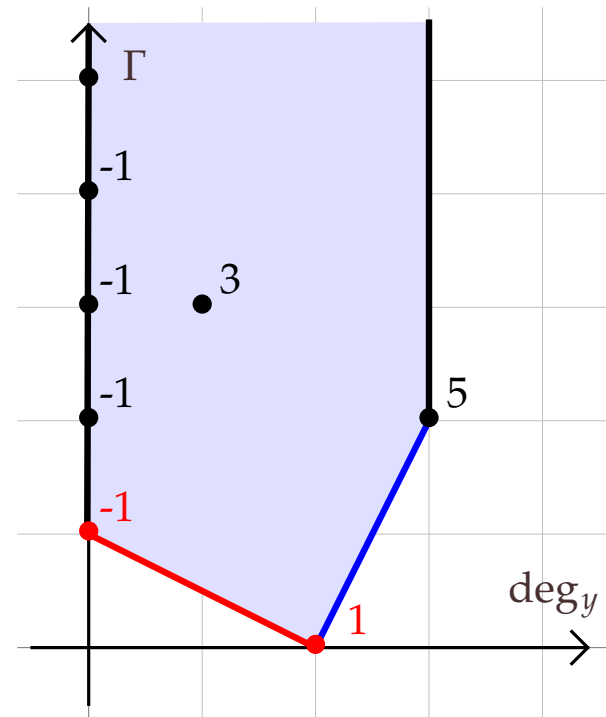


$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Assume $y \sim z^{1/2}$ and consider

$$y = z^{1/2} + \tilde{y} \quad (\tilde{y} < z^{1/2}).$$

Refinement := change of variable
 +
 asymptotic constraint



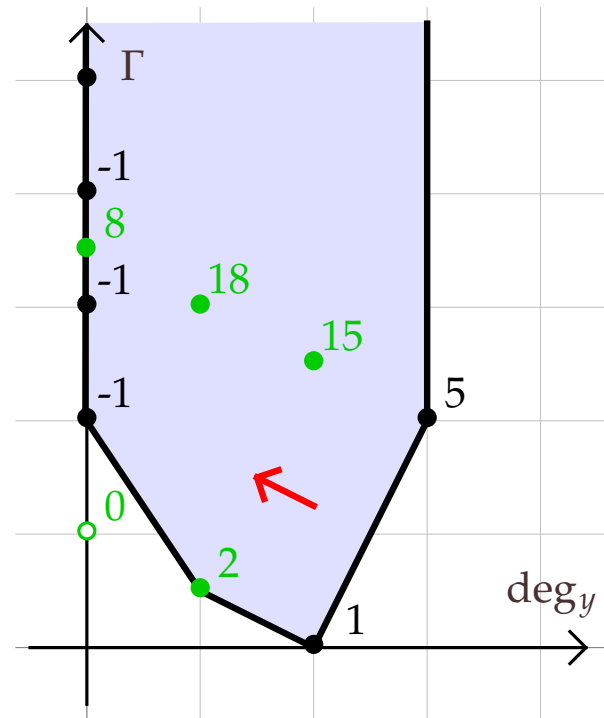
$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

The refinement

$$y = z^{1/2} + \tilde{y} \quad (\tilde{y} < z^{1/2}).$$

yields

$$\begin{aligned} & 5z^2\tilde{y}^3 \\ & + (1 + 15z^{2^{1/2}})\tilde{y}^2 \\ & + (2z^{1/2} + 18z^{2^{1/2}})\tilde{y} \\ & - z^2 - z^3 + 8z^{3^{1/2}} - z^4 - \dots = 0, \quad (\tilde{y} < z^{1/2}). \end{aligned}$$



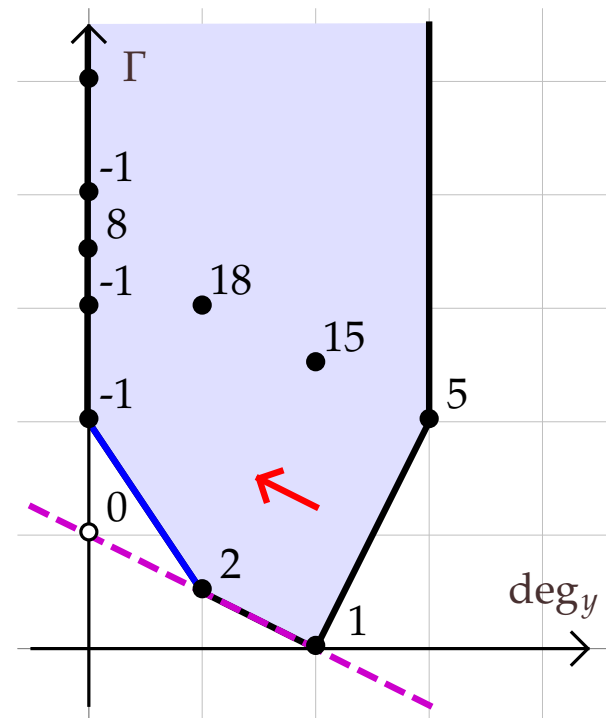
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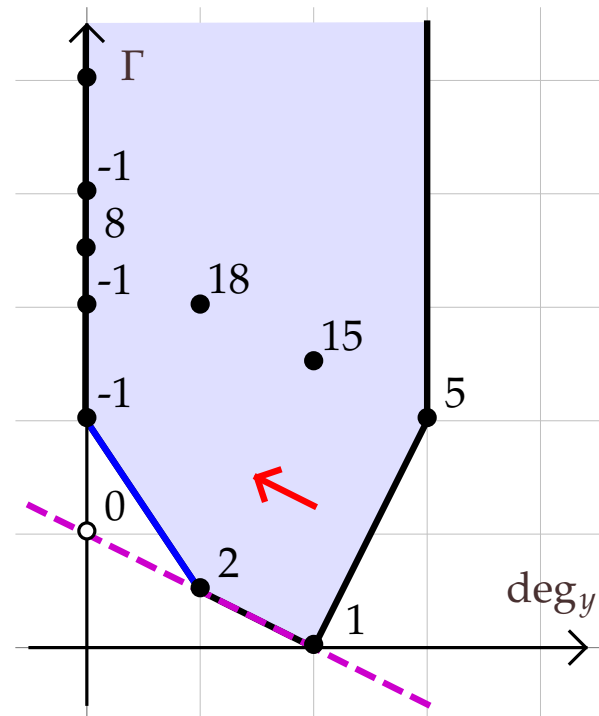
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Only new starting monomial: $\tilde{y} \asymp z^{3/2}$.

Only new starting monomial: $\tilde{y} \asymp \frac{1}{2}z^{3/2}$.



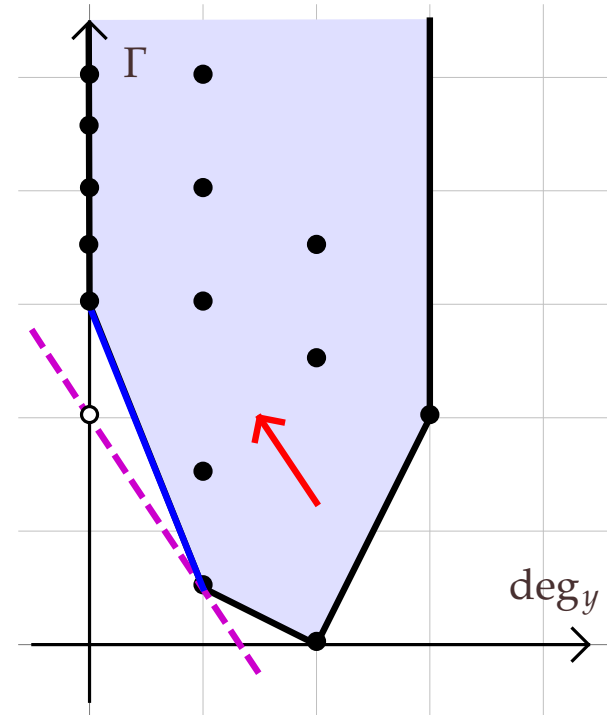
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Continued refinement process

$$y = z^{1/2} + \tilde{y} \quad (\tilde{y} < z^{1/2})$$

$$\tilde{y} = 1/2z^{3/2} + \tilde{\tilde{y}} \quad (\tilde{\tilde{y}} < z^{3/2})$$

⋮



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Continued refinement process

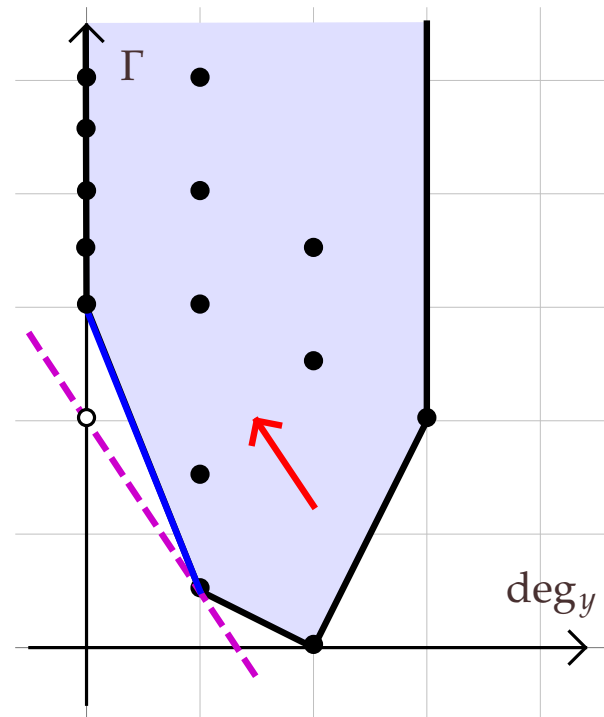
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⋮

yields asymptotic expansion

$$y \approx z^{1/2} + 1/2z^{3/2} + \dots$$



$$5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

Continued refinement process

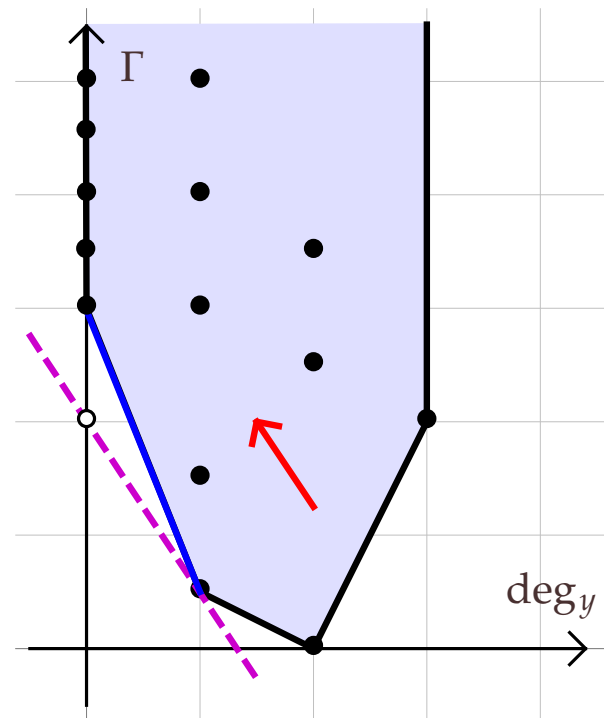
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⋮

yields asymptotic *solution*

$$y = z^{1/2} + 1/2z^{3/2} + \dots ?$$



$$P(y) = 5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z}$$

Multiplicative conjugate by $z^{1/2}$

$$\begin{aligned} P_{\times z^{1/2}}(y) &:= P(z^{1/2}y) \\ &= 5z^{3/2}y^3 + zy^2 + 3z^{3/2}y - \frac{z}{1-z} \end{aligned}$$

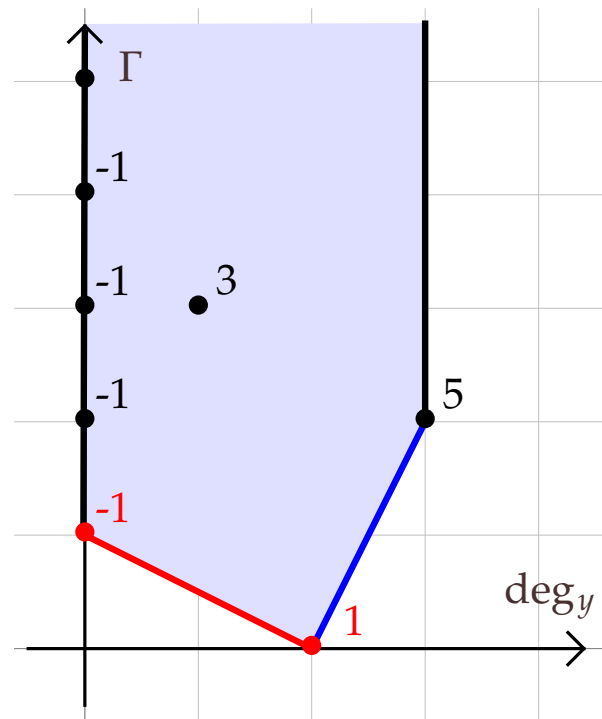
Multiplicative conjugations

$z^{1/2}$ is a starting monomial for

$$P(y) = 5z^2y^3 + y^2 + 3z^3y - \frac{z}{1-z} = 0$$

$\Leftrightarrow 1$ is a starting monomial for

$$P_{xz^{1/2}}(y) = 5z^{3^{1/2}}y^3 + zy^2 + 3z^{3^{1/2}}y - \frac{z}{1-z} = 0.$$



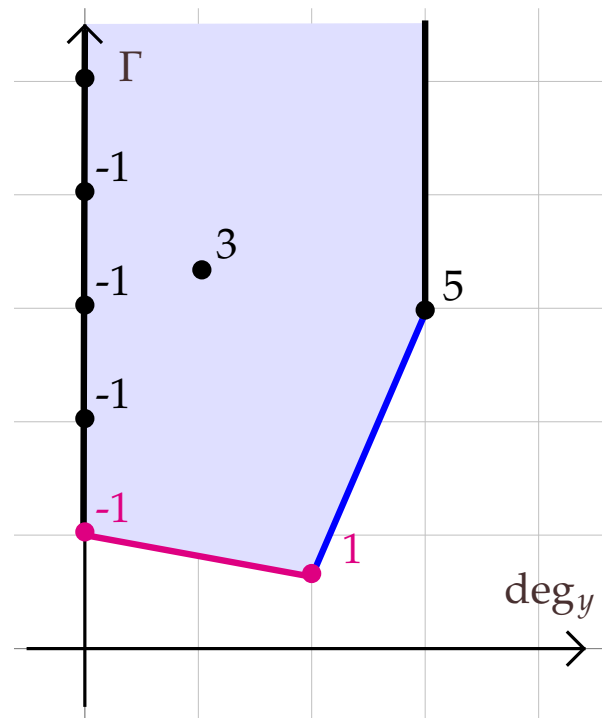
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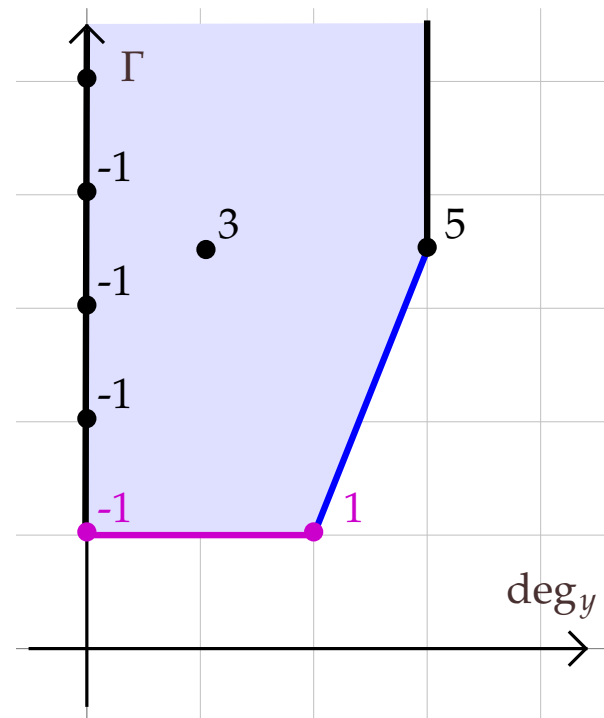


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Characterization of starting terms

cz^ν is a starting term for $P(y) = 0 \iff N_{P_{xz^\nu}}(c) = 0. \quad (c \neq 0)$

$$P(y) = 0 \quad (y < z^\gamma). \quad (\star)$$

Newton degree of (\star)

$$\deg_{<z^\gamma} P := \text{val } N_{P_{xz^\gamma}}$$

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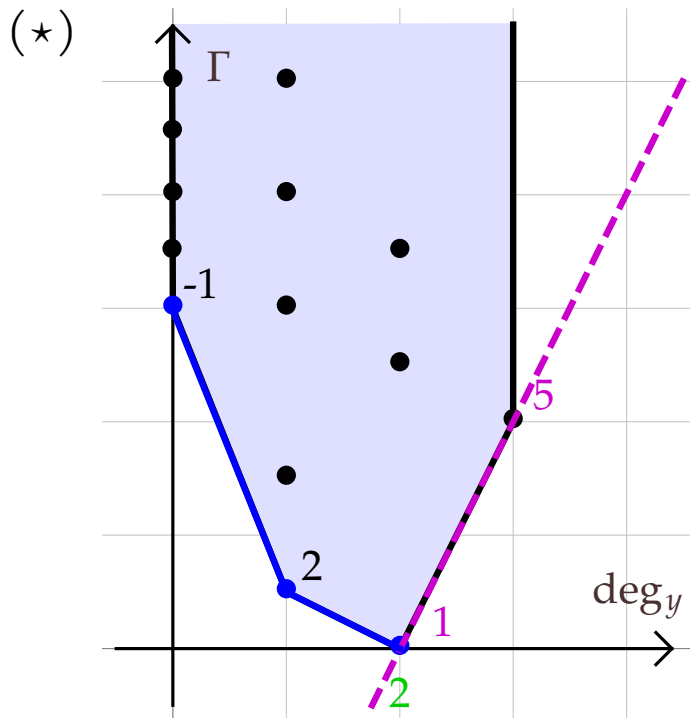
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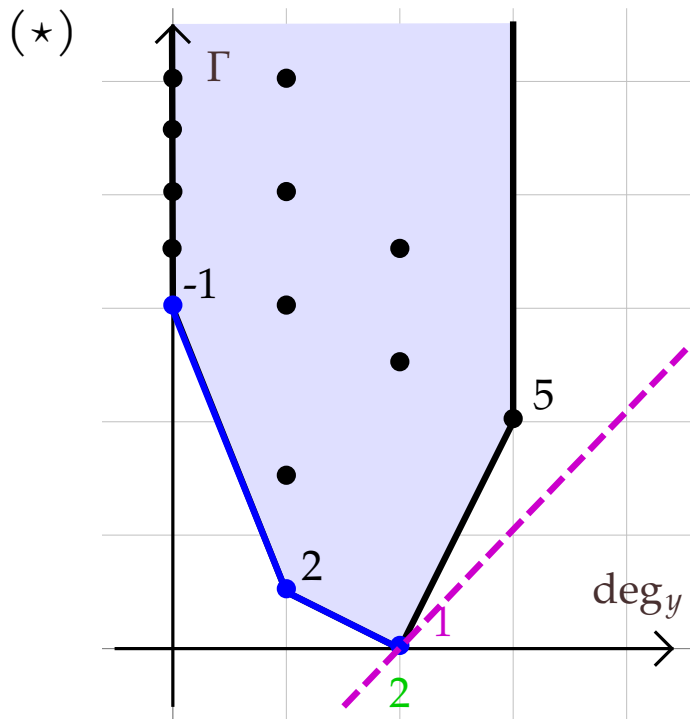
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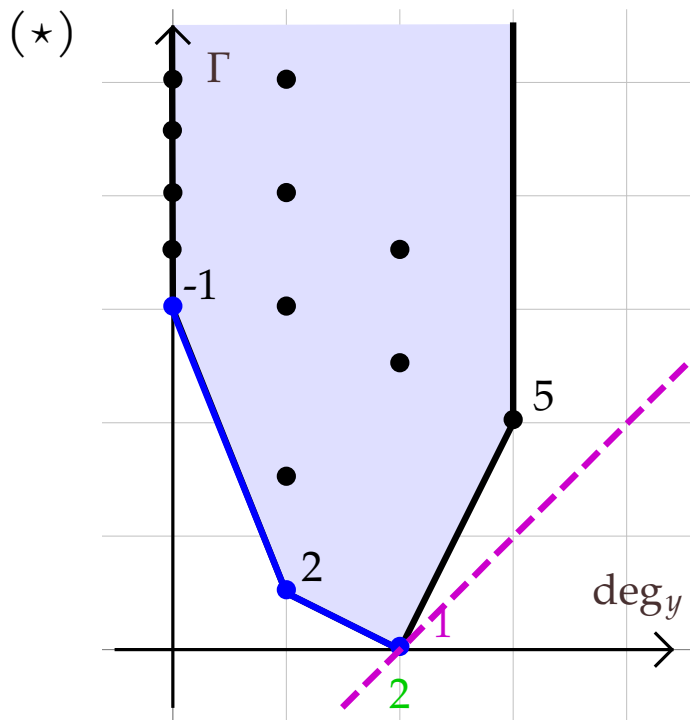
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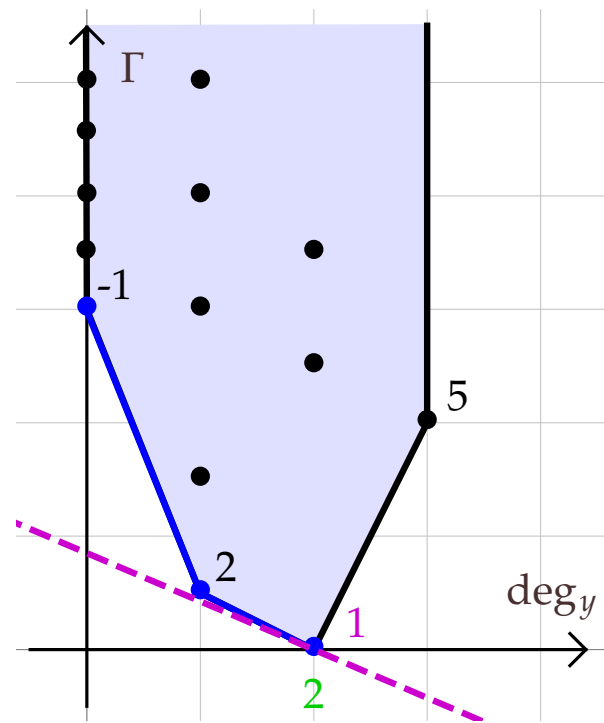
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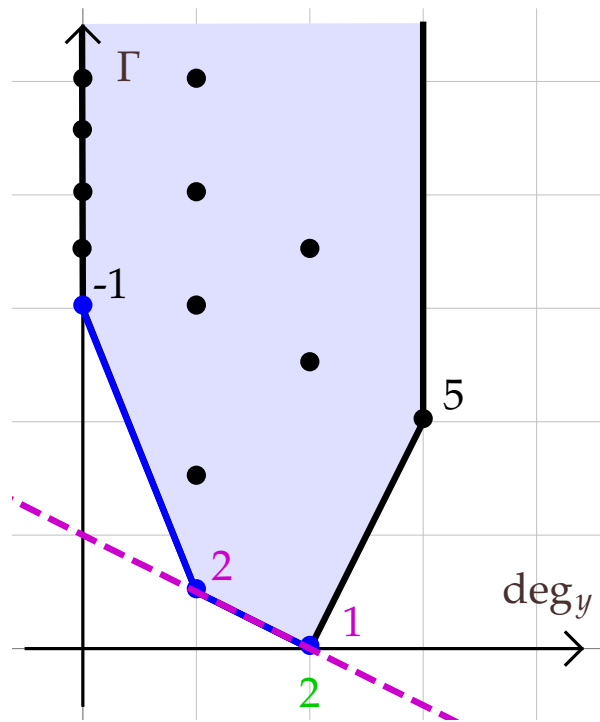
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Given $P \in K[[z^\Gamma]][Y]$ and $\varphi \in K[[z^\Gamma]]$, the **additive conjugate** of P by φ is

$$P_{+\varphi}(y) := P(\varphi + y)$$

Let $N \in K[Y]^{\neq 0}$ and let $c \in K$. Then

$$\text{val } N_{+c} = \text{multiplicity of } c \text{ as a root of } N$$

Let $c_1, \dots, c_\ell \in K$ be the roots of N . Since K is algebraically closed, we have

$$\deg N = \text{val } N_{+c_1} + \dots + \text{val } N_{+c_\ell}.$$

Consider an equation $P(y) = 0$, $y \prec z^\gamma$ of Newton degree d :

$$d = \deg_{\prec z^\gamma} P.$$

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If $\text{val}_Y P = d$, then $y = 0$ is a solution of multiplicity d .

Consider an equation $P(y) = 0$, $y \prec z^\gamma$ of Newton degree d :

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Assume that $\text{val}_\gamma P < d$ and let z^ν be the largest starting monomial. We have

$$d = \deg N, \quad N := N_{P_{\prec z^\nu}}.$$

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For any $\alpha \in K$, we have $P_{\times z^\nu, +\alpha} = P_{+\alpha z^\nu, \times z^\nu}$ and $N_{P_{+\alpha}} = N_{P, +\alpha}$, whence

$$\text{val } N_{+c_i} = \text{val } N_{P_{+c_i z^\nu, \times z^\nu}} = \deg_{<z^\nu} P_{+c_i z^\nu}.$$

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Hence

$$d = \deg_{\prec z^\nu} P_{+c_1 z^\nu} + \dots + \deg_{\prec z^\nu} P_{+c_\ell z^\nu}.$$

Conservation of Newton degree

Consider an asymptotic algebraic equation

$$P(y) = 0 \quad (y < z^\gamma),$$

with $\text{val}_\gamma P < \deg_{<z^\gamma} P$ and let z^ν be the largest starting monomial. Let c_1, \dots, c_ℓ be the roots of $N := N_{P_{<z^\nu}}$, so that each c_i determines a refined equation

$$P_{+c_i z^\nu}(\tilde{y}) = 0 \quad (\tilde{y} < z^\nu).$$

If K is algebraically closed, then

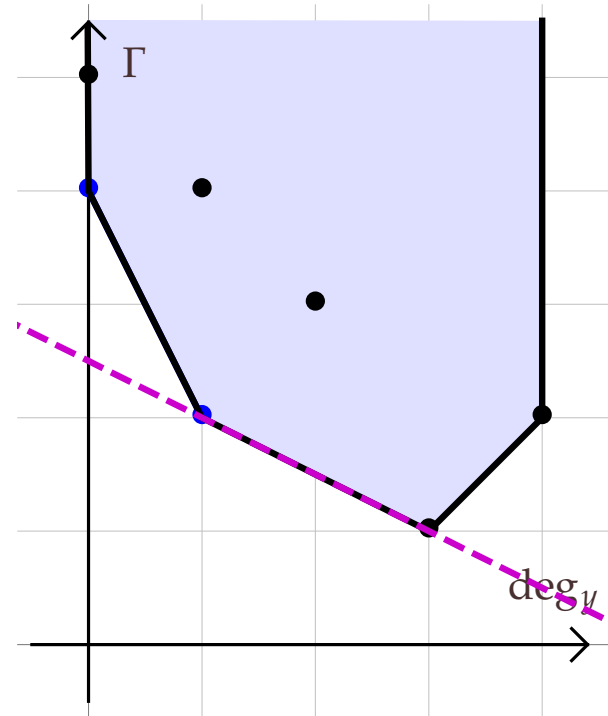
$$\deg_{<z^\gamma} P = \deg_{<z^\nu} P_{+c_1 z^\nu} + \dots + \deg_{<z^\nu} P_{+c_\ell z^\nu}.$$

Definition

The equation

$$P(y) = 0 \quad (y < z^\gamma)$$

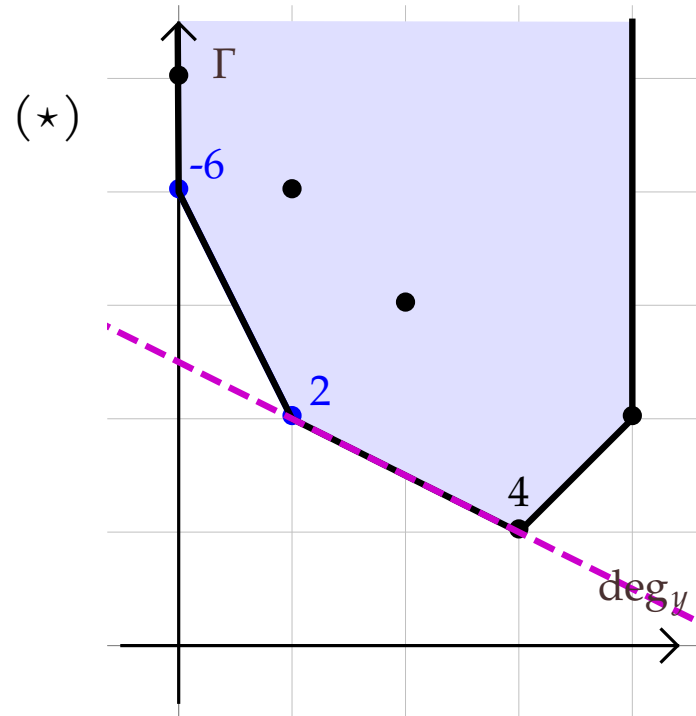
is *quasi-linear* if $\deg_{<z^\gamma} P = 1$.



Quasi-linear equations

Consider a quasi-linear equation

$$P(y) = 0 \quad (y < z^\gamma)$$

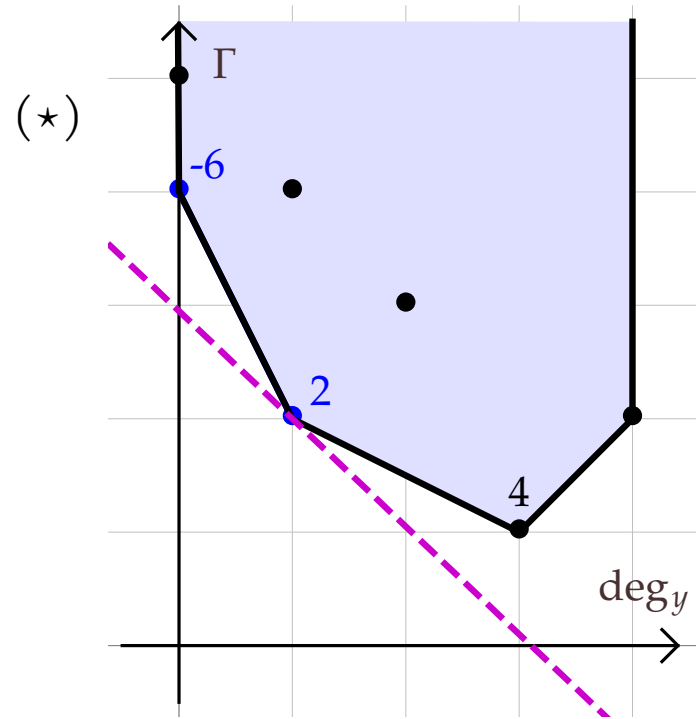


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Without loss of generality, we may arrange that

$$\text{val } N_{xz^\gamma} = \deg N_{xz^\gamma} = 1.$$

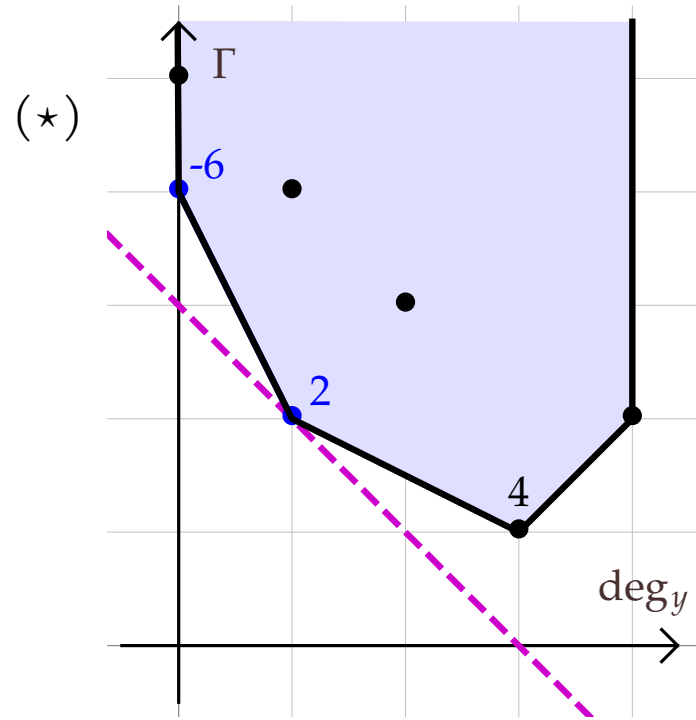


Quasi-linear equations

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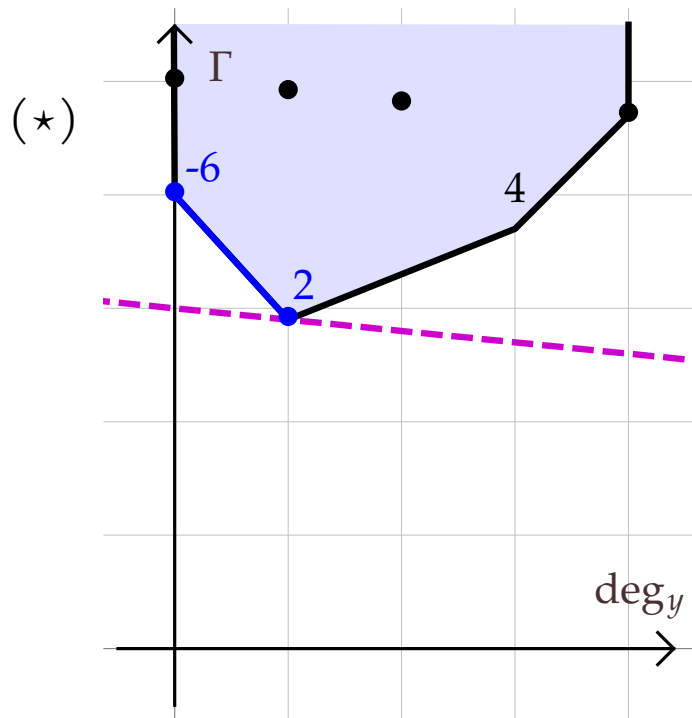
$$y := z^\gamma u$$

$$Q := P_{xz^\gamma}.$$

Then (\star) is equivalent to

$$Q(u) = 0 \quad (u < 1)$$

We have $\deg_{<1} Q = \text{val } N_Q = \deg N_Q = 1$.



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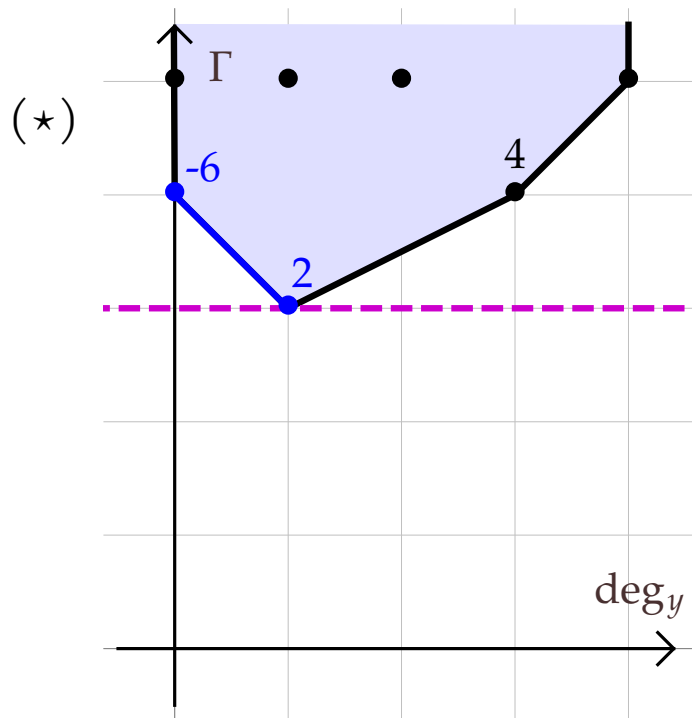
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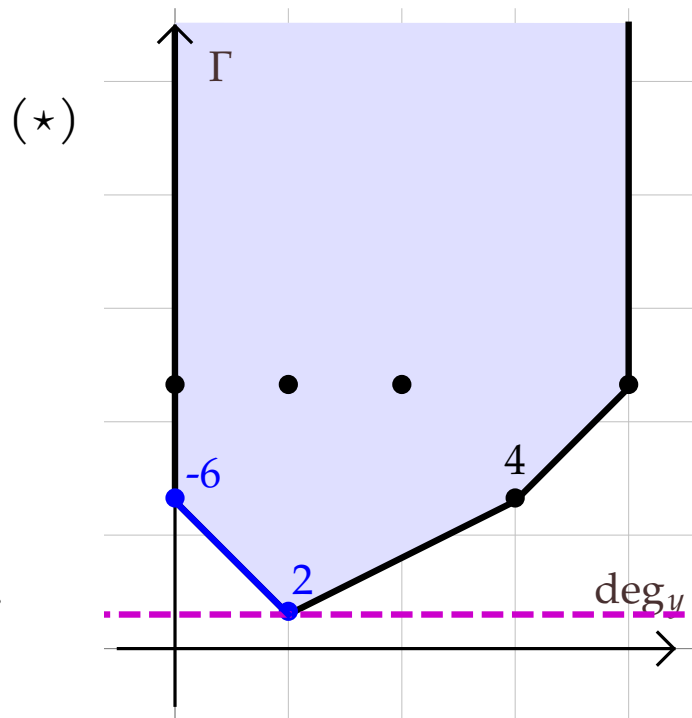
$$Q := P_{xz^\gamma}$$

$$R := \mathfrak{d}_Q^{-1} Q.$$

Then (\star) is equivalent to

$$R(u) = 0 \quad (u < 1)$$

We have $\deg_{<1} R = \text{val } N_R = \deg N_R = 1$ and $\mathfrak{d}_R = 1$.



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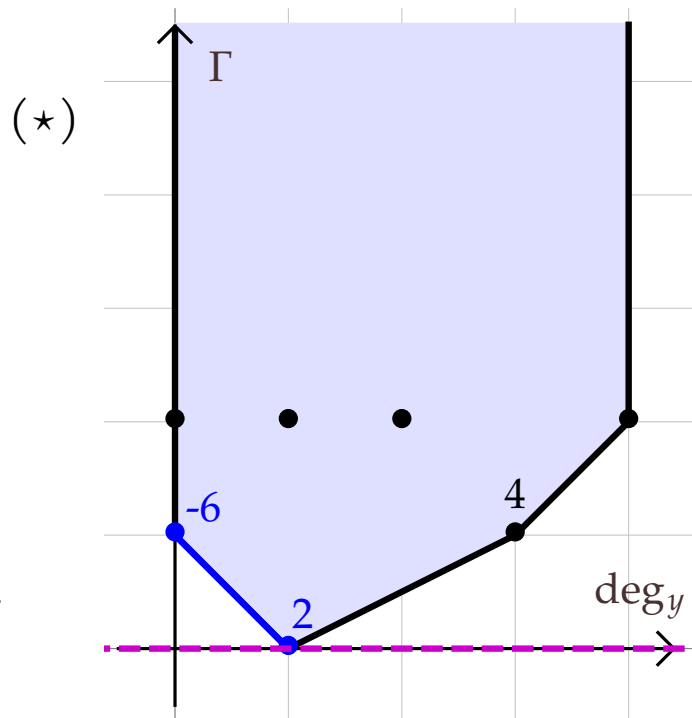
$$R := \vartheta_Q^{-1} Q.$$

Then (\star) is equivalent to

$$R(u) = 0 \quad (u < 1)$$

We have $\deg_{<1} R = \text{val } N_R = \deg N_R = 1$ and $\vartheta_R = 1$.

The polynomial R is in **Hensel position**.



Consider a quasi-linear equation

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Let

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$$Q := P_{\times z^\gamma}$$

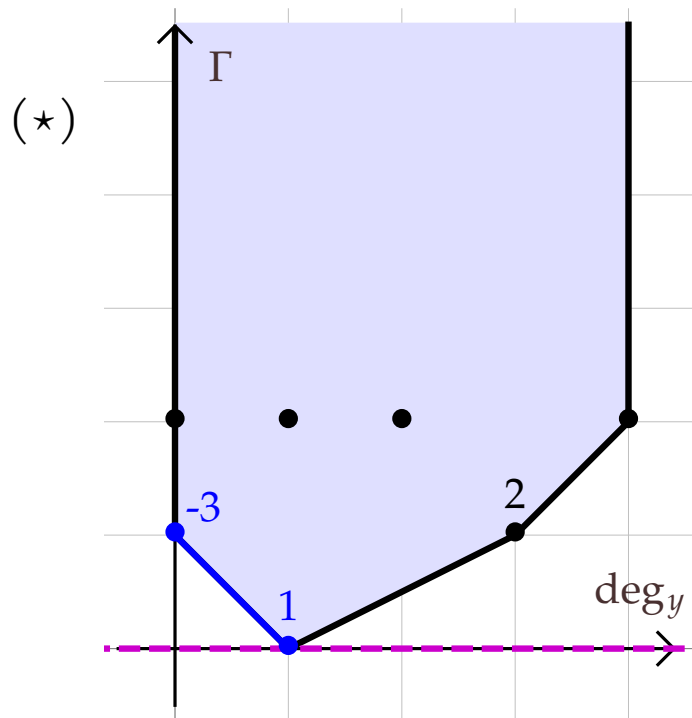
$$R := \partial_Q^{-1} Q$$

$$S := N_{R,1}^{-1} R.$$

Then $(*)$ is equivalent to

$$S(u) = 0 \quad (u < 1)$$

We have $\text{val } N_S = \text{deg } N_S = 1$, $\partial_S = 1$, and $N_{S,1} = 1$.



Consider a quasi-linear equation

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Let

$$y := z^\gamma u$$

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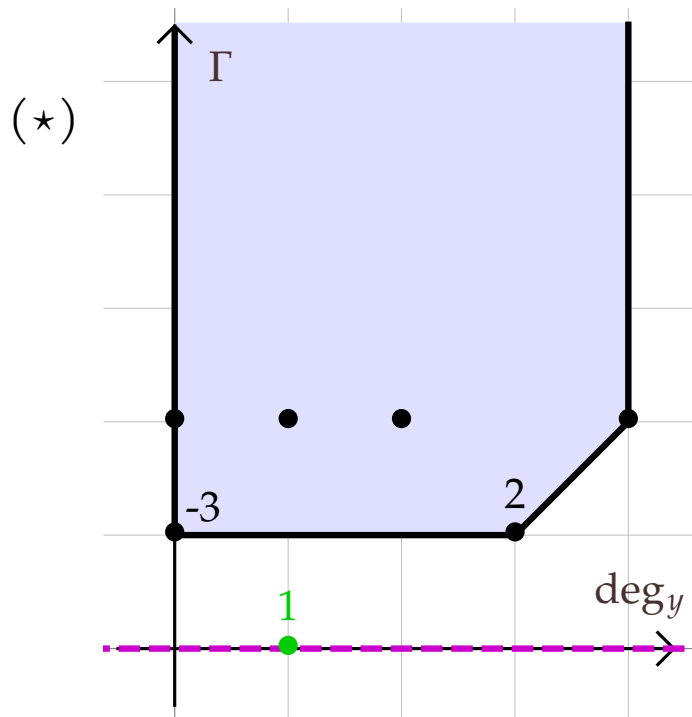
$$S := N_{R,1}^{-1} R$$

$$T := Y - S.$$

Then (\star) is equivalent to

$$u = T(u) \quad (u < 1)$$

We have $T < 1$. $u = -3z + 2zu^3 + O(z^2)$



Theorem

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Given $\varphi, \varepsilon < 1$ in $K[[z^\Gamma]]$, we have

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$$P(\varphi + \varepsilon) - P(\varphi) = P'(\varphi)\varepsilon + \frac{1}{2}P''(\varphi)\varepsilon^2 + \dots < \varepsilon.$$

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Given $n \in \mathfrak{S}$ with $n > m$, we have $y_{> m} - P(y_{> m}) = y_{\geq n} - P(y_{\geq n}) + o(n) < n$.

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Take $y_m := P(y_{> m})_m$.

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Let \mathfrak{M} be a monomial monoid and let $P \in K[[\mathfrak{M}]][Y]$ be such that $P < 1$ (i.e. $\text{supp } P < 1$). Then $y = P(y)$ has a unique solution in $K[[\mathfrak{M}]]^{<1}$.

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$$T = \begin{array}{c} \mathfrak{m} \\ \swarrow \quad \downarrow \quad \searrow \\ T_1 \quad \cdots \quad T_k \end{array} \quad \Rightarrow \quad \tau_T := P_{k,\mathfrak{m}} \mathfrak{m} \tau_{T_1} \cdots \tau_{T_k}.$$

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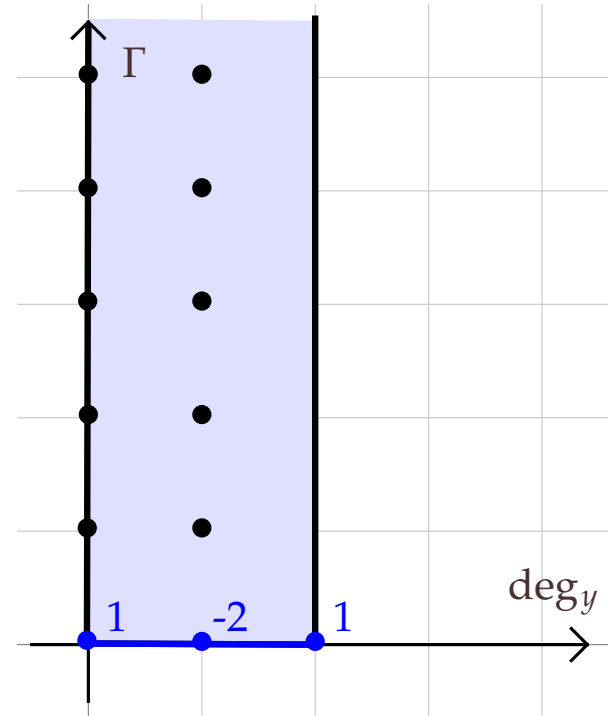
$$\begin{aligned} P(y) &= \sum_{k \in \mathbb{N}} \sum_{\mathfrak{m} \in \mathfrak{S}} P_{k,\mathfrak{m}} \mathfrak{m} y^k = \sum_{k \in \mathbb{N}} \sum_{\mathfrak{m} \in \mathfrak{S}} \sum_{T_1 \in \mathfrak{S}^\top} \cdots \sum_{T_k \in \mathfrak{S}^\top} P_{k,\mathfrak{m}} \mathfrak{m} \tau_{T_1} \cdots \tau_{T_k} \\ &= \sum_{T \in \mathfrak{S}^\top} \tau_T = y. \end{aligned}$$

□

Almost multiple solutions

Consider the equation

$$\left(y - \frac{1}{1-z}\right)^2 = z^{1000}.$$

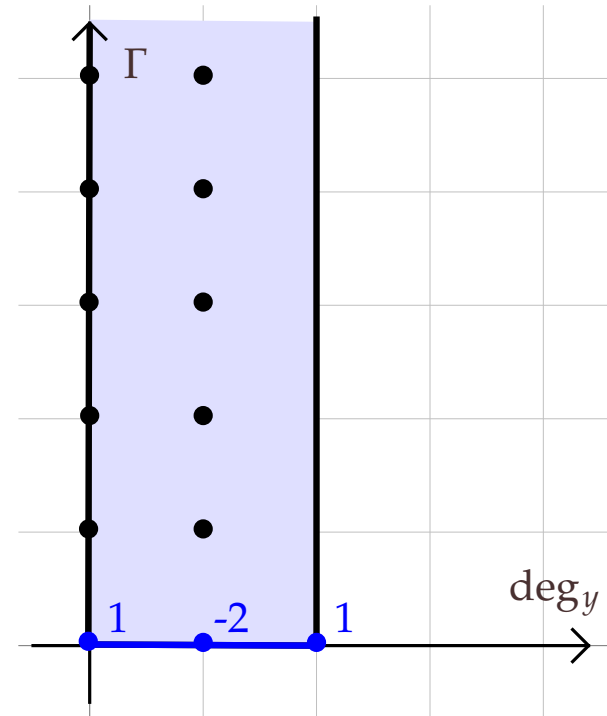


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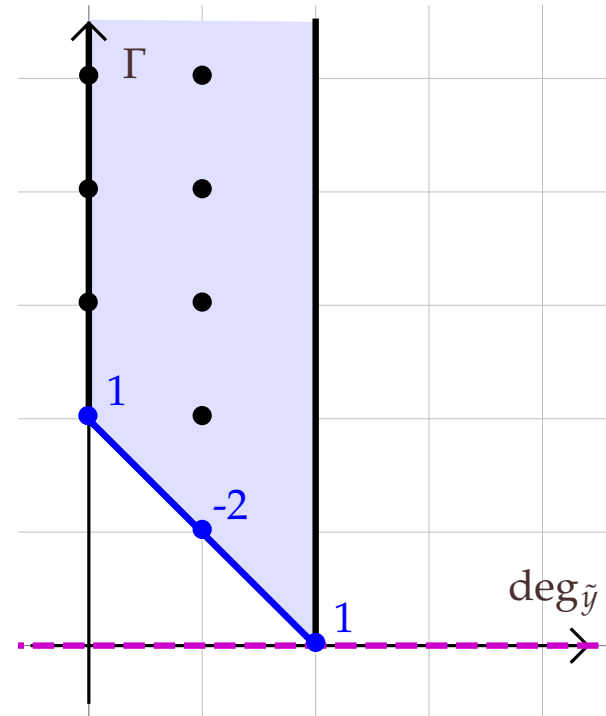
$$\left(y - \frac{1}{1-z}\right)^2 = z^{1000}.$$

There is a unique starting term $y \sim 1$. After

$$y = 1 + \tilde{y} \quad (\tilde{y} < 1),$$

we obtain

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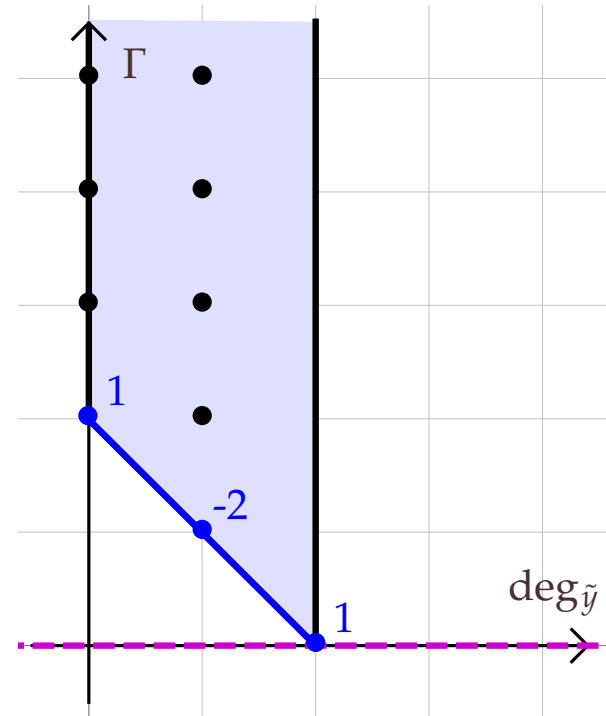
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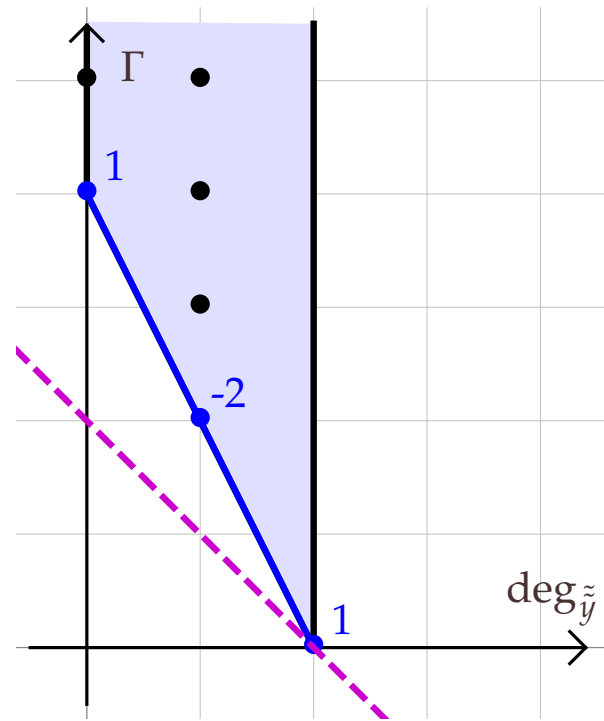
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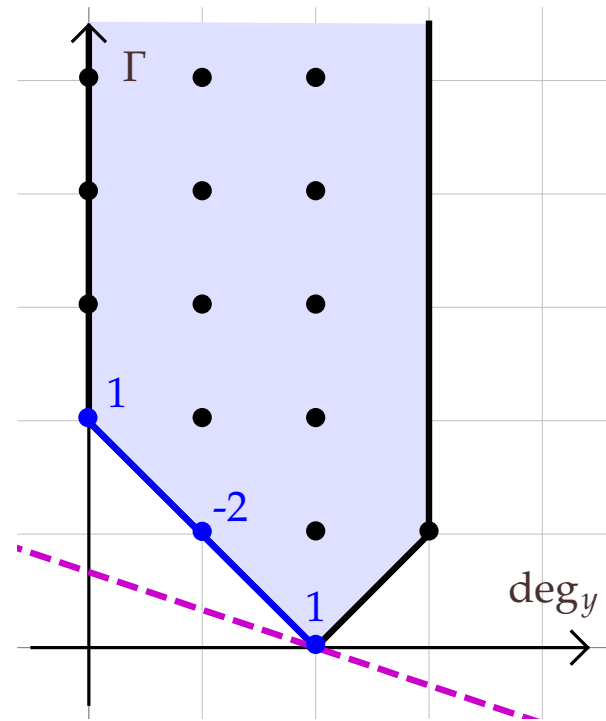
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$$P(y) = 0 \quad (y < z^\gamma)$$

with unique d -fold starting term $y \sim cz^\nu$.

Then $N_{P_{xz^\nu}} = \alpha(Y - c)^d$, where $d = \deg_{<z^\gamma} P$.

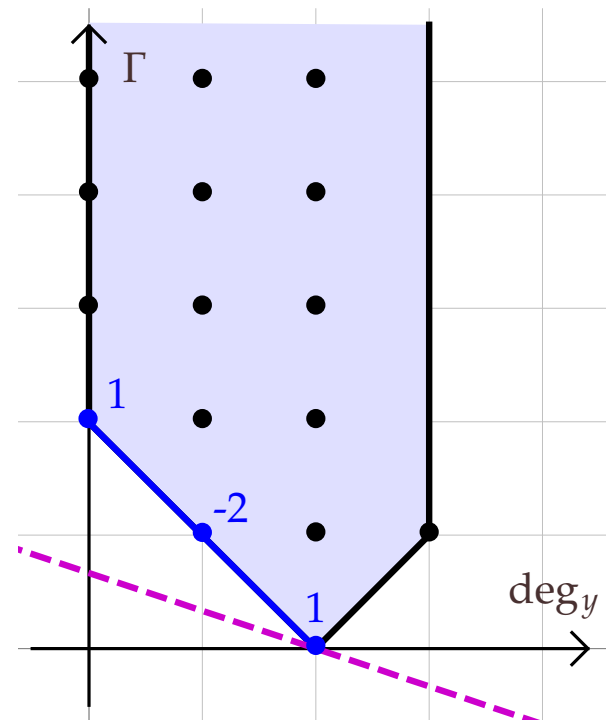


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Note: $\text{char } K = 0 \implies N_{P_{xz^\nu}, d-1} = -dc \neq 0 \implies P_{d-1} \neq 0$.



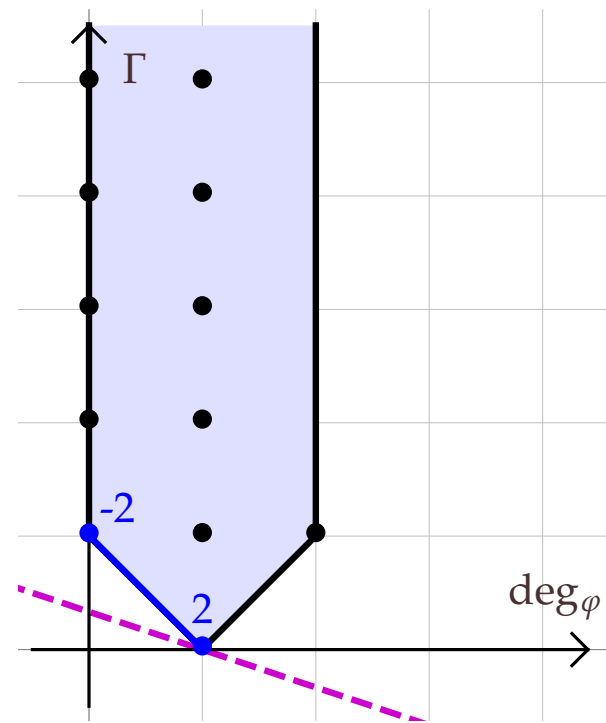
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φ : the unique solution of the quasi-linear equation

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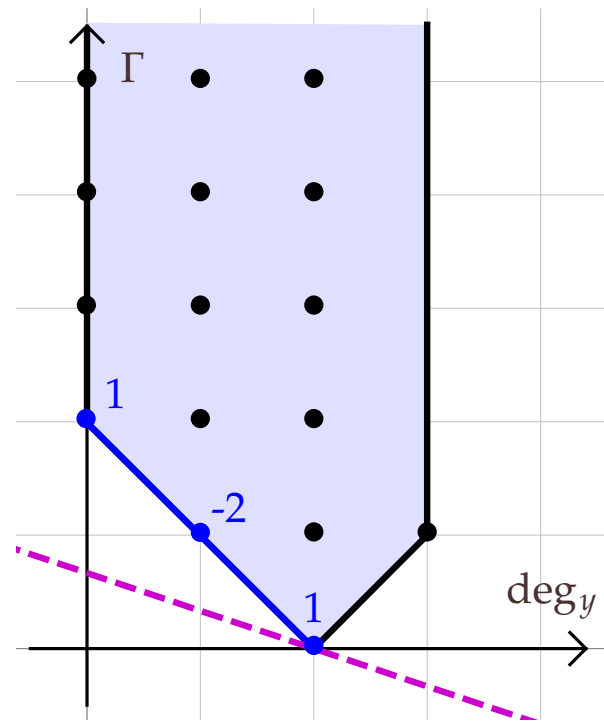
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instead of

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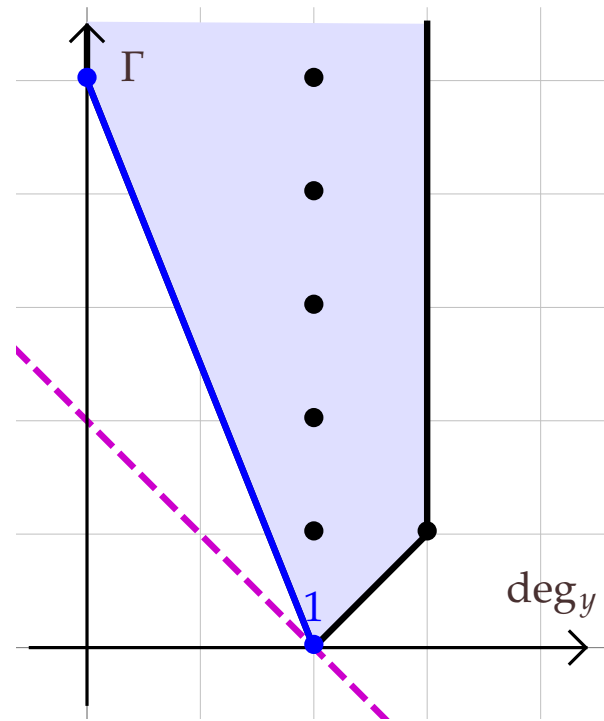
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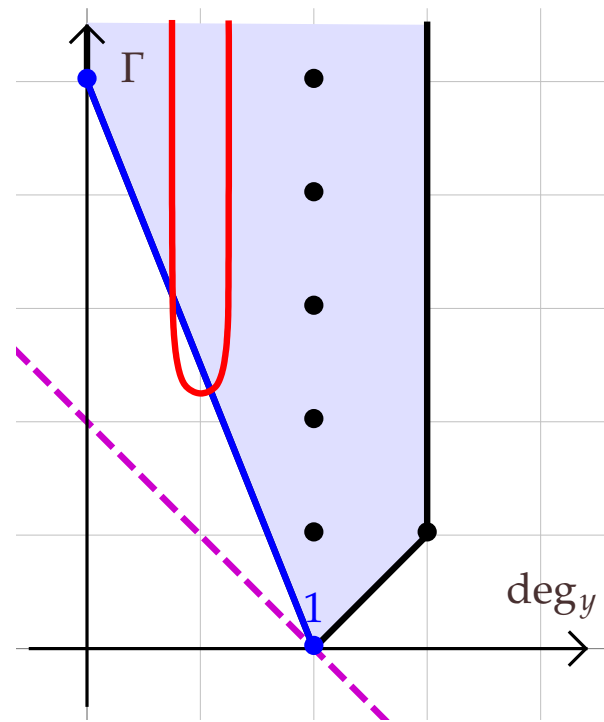
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Then

$$P_{+\varphi}(\tilde{y}) = P(\varphi) + P'(\varphi)\tilde{y} + \dots = 0 \quad (\tilde{y} < z^\nu),$$

whence $P_{+\varphi, d-1} = 0$



Algorithm solve(P, z^γ)

INPUT: $P \in K[[z^\Gamma]][Y]$ and $z^\gamma \in z^\Gamma$ with $d := \deg_{<z^\gamma} P > 0$ and $\text{char } K = 0$

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 - Compute $\varphi := \text{solve}\left(\frac{\partial^{d-1} P}{\partial Y^{d-1}}, z^\gamma\right)$
 - Compute $\tilde{y}_1, \dots, \tilde{y}_d := \text{solve}(P_{+\varphi}, z^\nu)$
 - Return $\varphi + \tilde{y}_1, \dots, \varphi + \tilde{y}_d$
- For $i = 1, \dots, \ell$, compute $y_{i,1}, \dots, y_{i,d_i} := \text{solve}(P_{+c_i z^\nu}, z^\nu)$

Algorithm solve(P, z^γ)

INPUT: $P \in K[[z^\Gamma]][Y]$ and $z^\gamma \in z^\Gamma$ with $d := \deg_{<z^\gamma} P > 0$ and $\text{char } K = 0$

OUTPUT: solutions $y_1, \dots, y_d \in K[[z^\Gamma]]$ of $P(y) = 0, y < z^\gamma$, counted with multiplicities

- If $\text{val}_Y P = d$, then return $0, \dots, 0$.
- If $d = 1$, then $P(y) = 0, y < z^\gamma$ is quasi-linear; return its unique solution
- Let z^ν be the largest starting monomial and let c_1, \dots, c_ℓ be the roots of $N_{P_{z^\nu}}$
- If $\ell = 1$, then
 - Compute $\varphi := \text{solve}\left(\frac{\partial^{d-1} P}{\partial Y^{d-1}}, z^\gamma\right)$
 - Compute $\tilde{y}_1, \dots, \tilde{y}_d := \text{solve}(P_{+\varphi}, z^\nu)$
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- For $i = 1, \dots, \ell$, compute $y_{i,1}, \dots, y_{i,d_i} := \text{solve}(P_{+c_i z^\nu}, z^\nu)$
- Return $y_{1,1}, \dots, y_{1,d_1}, \dots, y_{\ell,1}, \dots, y_{\ell,d_\ell}$

Let K be a field of characteristic zero.

Theorem

Let $P \in K[[z^\Gamma]][Y]^{\neq 0}$ and $z^\gamma \in z^\Gamma$. If K is algebraically closed and Γ divisible, then

$$P(y) = 0 \quad (y < z^\gamma)$$

has exactly $\deg_{< z^\gamma} P$ solutions in $K[[z^\Gamma]]$, when counting with multiplicities.

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- Generalizations to $K[[z^\Gamma]]_{\mathcal{F}}$ and $K[[z^\Gamma]]_{\mathcal{L}}$ instead of $K[[z^\Gamma]]$.