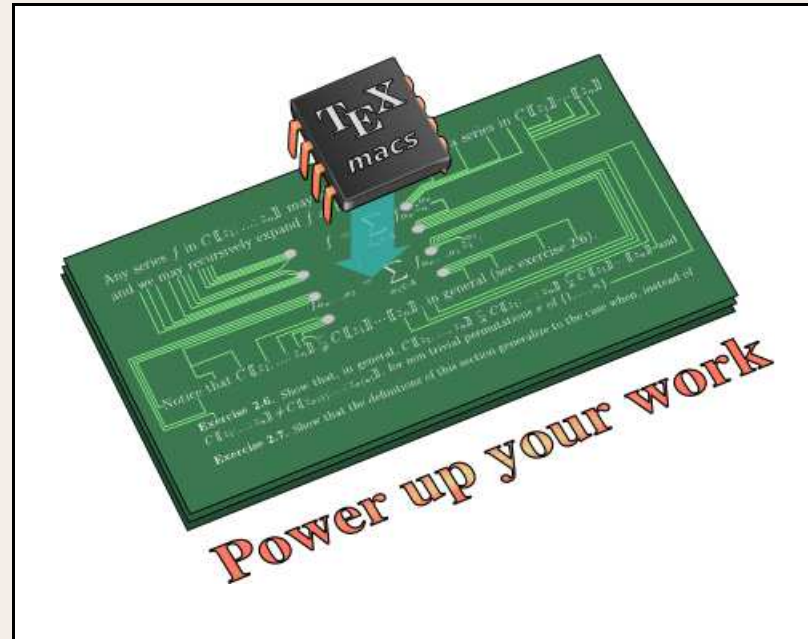


# Asymptotic differential equations

## Lecture 1: transseries and asymptotic differential equations



Joris van der Hoeven, Segovia 2011

<http://www.TEXMACS.org>



# Asymptotic differential algebra



Algebraic geometry



Real algebraic geometry  
+  
Valuation theory



Differential algebra



?

- Hardy fields: Rosenlicht, Boshernitzan, Singer, etc.
- Pfaff systems: Khovanskii, Wilkie, etc.



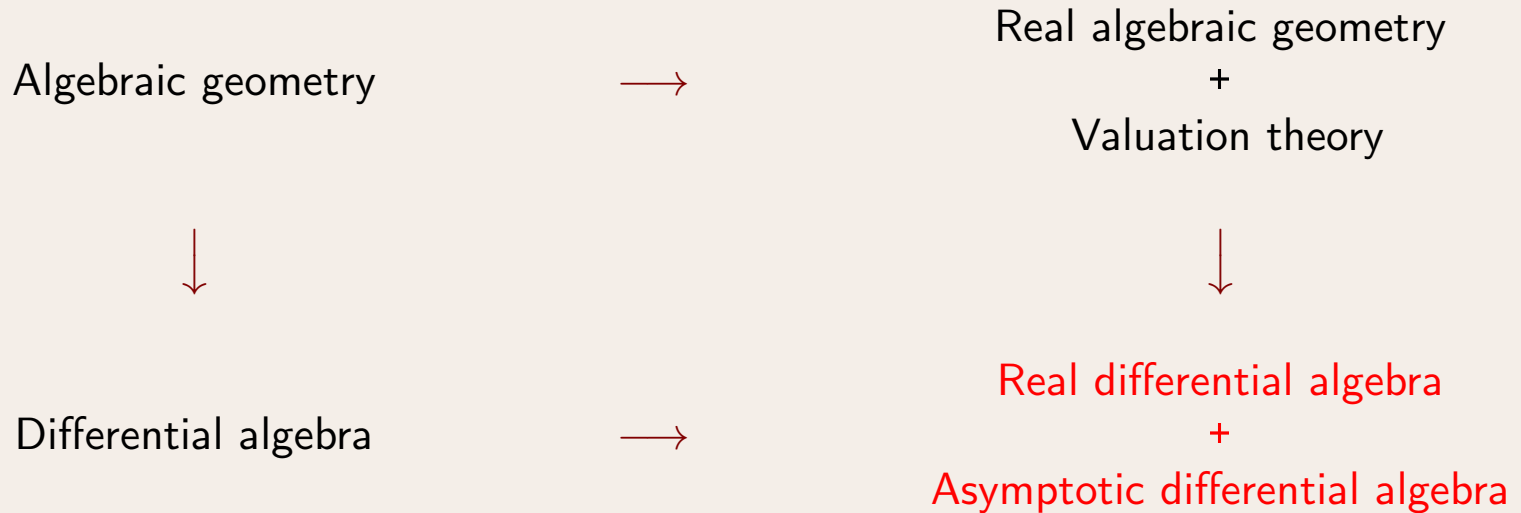
# Asymptotic differential algebra



- Hardy fields: Rosenlicht, Boshernitzan, Singer, etc.
- Pfaff systems: Khovanskii, Wilkie, etc.



# Asymptotic differential algebra



- LNM 1888: Transseries and Real Differential Algebra
- Upcoming book with Matthias Aschenbrenner and Lou van den Dries



# Sufficiently closed models



Algebraic geometry



Real algebraic geometry  
+  
Valuation theory



Differential algebra



Real differential algebra  
+  
Asymptotic differential algebra



# Sufficiently closed models



$\mathbb{C}$



Real algebraic geometry  
+  
Valuation theory



Differential algebra



Real differential algebra  
+  
Asymptotic differential algebra



# Sufficiently closed models



$\mathbb{C}$



$\mathbb{R}$

+

Valuation theory



Differential algebra



Real differential algebra

+

Asymptotic differential algebra



# Sufficiently closed models

 $\mathbb{C}$  $\longrightarrow$  $\mathbb{R}$   
+  
 $\mathbb{C}[[z^{\mathbb{Q}}]]$  $\downarrow$  $\downarrow$ 

Differential algebra

 $\longrightarrow$ Real differential algebra  
+  
Asymptotic differential algebra





# Sufficiently closed models



$\mathbb{C}$



$\mathbb{R}$   
+  
 $\mathbb{C}[[z^{\mathbb{Q}}]]$



*Wild*



Real differential algebra  
+  
Asymptotic differential algebra



# Sufficiently closed models



$\mathbb{C}$

$\longrightarrow$

$\mathbb{R}$   
+  
 $\mathbb{C}[[z^{\mathbb{Q}}]]$

$\downarrow$

$\downarrow$

*Wild*

$\longrightarrow$

Maximal Hardy field  
+  
Asymptotic differential algebra



# Sufficiently closed models



$\mathbb{C}$

$\longrightarrow$

$\mathbb{R}$   
+  
 $\mathbb{C} \llbracket z^{\mathbb{Q}} \rrbracket$

$\downarrow$

$\downarrow$

*Wild*

$\longrightarrow$

$\mathbb{R} \llbracket x \rrbracket$   
+

Asymptotic differential algebra



# Sufficiently closed models



$\mathbb{C}$

$\longrightarrow$

$\mathbb{R}$   
+  
 $\mathbb{C} \llbracket z^{\mathbb{Q}} \rrbracket$

$\downarrow$

$\downarrow$

*Wild*

$\longrightarrow$

$\mathbb{R} \llbracket x \rrbracket$   
+  
 $\mathbb{C} \llbracket z \rrbracket$



# Sufficiently closed models



$(x \succ 1)$

$$e^{e^x + \frac{e^x}{x} + \frac{e^x}{x^2} + \dots} + \frac{2}{\log x} e^{e^x + \frac{e^x}{x} + \frac{e^x}{x^2} + \dots} + e^{\sqrt{x} + e^{\sqrt{\log x}} + e^{\sqrt{\log \log x}} + \dots} + \dots$$



# Sufficiently closed models



$(x \succ 1)$

$$e^{e^x + \frac{e^x}{x} + \frac{e^x}{x^2} + \dots} + \frac{2}{\log x} e^{e^x + \frac{e^x}{x} + \frac{e^x}{x^2} + \dots} + e^{\sqrt{x} + e^{\sqrt{\log x}} + e^{\sqrt{\log \log x}} + \dots} + \dots$$

- Dahn & Göring
- Écalle



# Examples of transseries



$$\frac{1}{1 - x^{-1} - x^{-e}} = 1 + x^{-1} + x^{-2} + x^{-e} + x^{-3} + x^{-e-1} + \dots$$

$$\frac{1}{1 - x^{-1} + e^{-x}} = 1 + \frac{1}{x} + \frac{1}{x^2} + \dots + e^{-x} + 2 \frac{e^{-x}}{x} + \dots + e^{-2x} + \dots$$

$$-e^x \int \frac{e^{-x}}{x} = \frac{1}{x} - \frac{1}{x^2} + \frac{2}{x^3} - \frac{6}{x^4} + \frac{24}{x^5} - \frac{120}{x^6} + \dots$$

$$\Gamma(x) = \frac{\sqrt{2\pi} e^{x(\log x - 1)}}{x^{1/2}} + \frac{\sqrt{2\pi} e^{x(\log x - 1)}}{12 x^{3/2}} + \frac{\sqrt{2\pi} e^{x(\log x - 1)}}{288 x^{5/2}} + \dots$$

$$\zeta(x) = 1 + 2^{-x} + 3^{-x} + 4^{-x} + \dots$$

$$\varphi(x) = \frac{1}{x} + \varphi(x^\pi) = \frac{1}{x} + \frac{1}{x^\pi} + \frac{1}{x^{\pi^2}} + \frac{1}{x^{\pi^3}} + \dots$$

$$\psi(x) = \frac{1}{x} + \psi(e^{\log^2 x}) = \frac{1}{x} + \frac{1}{e^{\log^2 x}} + \frac{1}{e^{\log^4 x}} + \frac{1}{e^{\log^8 x}} + \dots$$



# MATHEMAGIX examples



```
Mmx] use "asymptotix"
```

```
Mmx] x == infinity ('x);
```

```
Mmx] 1 / (x + 1)
```

$$\frac{1}{x} - \frac{1}{x^2} + \frac{1}{x^3} - \frac{1}{x^4} + O\left(\frac{1}{x^5}\right)$$



**Mmx]** 1 / (exp x + x + 1)

$$\frac{1}{e^x} - \frac{x}{e^{2x}} - \frac{1}{e^{2x}} + \frac{x^2}{e^{3x}} + \frac{2x}{e^{3x}} + \frac{1}{e^{3x}} - \frac{x^3}{e^{4x}} - \frac{3x^2}{e^{4x}} - \frac{3x}{e^{4x}} - \frac{1}{e^{4x}} + O\left(\frac{x^4}{e^{5x}}\right)$$

**Mmx]** lengthen (exp (x^4 / (x + 1)), 7)

$$\frac{e^{x^3-x^2+x}}{e} + \frac{e^{x^3-x^2+x}}{ex} - \frac{e^{x^3-x^2+x}}{2ex^2} + O\left(\frac{e^{x^3-x^2+x}}{x^7}\right)$$

**Mmx]** integrate (exp (x^2), x)

$$\frac{e^{x^2}}{2x} + \frac{e^{x^2}}{4x^3} + \frac{3e^{x^2}}{8x^5} + \frac{15e^{x^2}}{16x^7} + O\left(\frac{e^{x^2}}{x^9}\right)$$

**Mmx]** integrate (x^x, x)

$$\frac{e^{x \log(x)}}{\log(x)} - \frac{e^{x \log(x)}}{\log(x)^2} + \frac{e^{x \log(x)}}{\log(x)^3} - \frac{e^{x \log(x)}}{\log(x)^4} + O\left(\frac{e^{x \log(x)}}{\log(x)^5}\right) + \frac{e^{x \log(x)}}{x \log(x)^3} - \frac{3e^{x \log(x)}}{x \log(x)^4} + \frac{6e^{x \log(x)}}{x \log(x)^5} + O\left(\frac{e^{x \log(x)}}{x \log(x)^6}\right) + \frac{e^{x \log(x)}}{x^2 \log(x)^4} - \frac{e^{x \log(x)}}{x^2 \log(x)^5} + O\left(\frac{e^{x \log(x)}}{x^2 \log(x)^6}\right) + \frac{2e^{x \log(x)}}{x^3 \log(x)^5} + O\left(\frac{e^{x \log(x)}}{x^4 \log(x)^6}\right)$$

**Mmx]** fixed\_point\_transseries (f :-> 1/x + f @ (x^2) + f @ (exp x))

$$\frac{1}{x} + \frac{1}{x^2} + \frac{1}{x^4} + \frac{1}{x^8} + O\left(\frac{1}{x^{16}}\right) + \frac{1}{e^x} + \frac{1}{e^{2x}} + \frac{1}{e^{4x}} + O\left(\frac{1}{e^{8x}}\right) + \frac{1}{e^{x^2}} + \frac{1}{e^{2x^2}} + O\left(\frac{1}{e^{4x^2}}\right) + \frac{1}{e^{x^4}} + O\left(\frac{1}{e^{2x^4}}\right) + \frac{1}{e^{e^x}} + \frac{1}{e^{2e^x}} + O\left(\frac{1}{e^{4e^x}}\right) + \frac{1}{e^{e^{2x}}} + O\left(\frac{1}{e^{2e^{2x}}}\right) + \frac{1}{e^{e^{x^2}}} + O\left(\frac{1}{e^{2e^{x^2}}}\right) + \frac{1}{e^{e^{e^x}}} + O\left(\frac{1}{e^{2e^{e^x}}}\right)$$

**Mmx]** lengthen (product (x, x), 4)

$$\frac{e^{x \log(x) - x}}{\text{sqrt}(x)} + \frac{e^{x \log(x) - x}}{12 x^{\frac{3}{2}}} + \frac{e^{x \log(x) - x}}{288 x^{\frac{5}{2}}} - \frac{139 e^{x \log(x) - x}}{51840 x^{\frac{7}{2}}} + O\left(\frac{e^{x \log(x) - x}}{x^{\frac{9}{2}}}\right)$$

**Mmx]** lengthen (product (log x, x), 2)

$$\begin{aligned} & \frac{e^{x \log(\log(x)) - \frac{x}{\log(x)} - \frac{x}{\log(x)^2} - \frac{2x}{\log(x)^3} + O\left(\frac{x}{\log(x)^4}\right) - \frac{\log(\log(x))}{2}}{12 x \log(x)} + \\ & \frac{e^{x \log(\log(x)) - \frac{x}{\log(x)} - \frac{x}{\log(x)^2} - \frac{2x}{\log(x)^3} + O\left(\frac{x}{\log(x)^4}\right) - \frac{\log(\log(x))}{2}}{288 x^2 \log(x)^2} - \\ & \frac{e^{x \log(\log(x)) - \frac{x}{\log(x)} - \frac{x}{\log(x)^2} - \frac{2x}{\log(x)^3} + O\left(\frac{x}{\log(x)^4}\right) - \frac{\log(\log(x))}{2}}{360 \log(x) x^3} - \\ & \frac{e^{x \log(\log(x)) - \frac{x}{\log(x)} - \frac{x}{\log(x)^2} - \frac{2x}{\log(x)^3} + O\left(\frac{x}{\log(x)^4}\right) - \frac{\log(\log(x))}{2}}{240 x^3 \log(x)^2} + \\ & O\left(\frac{e^{x \log(\log(x)) - \frac{x}{\log(x)} - \frac{x}{\log(x)^2} - \frac{2x}{\log(x)^3} + O\left(\frac{x}{\log(x)^4}\right) - \frac{\log(\log(x))}{2}}{x^3 \log(x)^3}\right) \end{aligned}$$

**Mmx]** eval (integrate (exp (x^2), x), x, 100.0)

4.40362931632092710468e4340

**Mmx]**



# Generalized power series



- $C$ : constant field
- $(\mathfrak{M}, \preceq)$ : totally ordered group of monomials  
 i.e.  $\log \mathfrak{M}$  is a value group with  $\mathfrak{m} \preceq \mathfrak{n} \Leftrightarrow v(\log \mathfrak{m}) \geq v(\log \mathfrak{n})$
- $C[[\mathfrak{M}]]$ : Hahn field of  $f = \sum_{\mathfrak{m} \in \mathfrak{M}} f_{\mathfrak{m}} \mathfrak{m}$  with **well-based support**.

$\mathfrak{m}_1 \prec \mathfrak{m}_2 \prec \dots$  with  $\mathfrak{m}_1, \mathfrak{m}_2, \dots \in \text{supp } f$  is impossible

- $C[[\mathfrak{M}]]$ : field of  $f = \sum_{\mathfrak{m} \in \mathfrak{M}} f_{\mathfrak{m}} \mathfrak{m}$  with **grid-based support**.

$$\text{supp } f \subseteq \mathfrak{m}_1^{\mathbb{N}} \cdots \mathfrak{m}_m^{\mathbb{N}} \mathfrak{n}, \quad \mathfrak{m}_1, \dots, \mathfrak{m}_m \prec 1$$

- $\mathcal{S} \subseteq \mathcal{P}(\mathfrak{M})$  closed under  $\cup, \cdot$  and power products of infinitesimal sets, with  $\{\mathfrak{m}\} \in \mathcal{S}$  for all  $\mathfrak{m} \in \mathfrak{M}$ .



# Generalized power series



- $C$ : constant ring (or set)
- $(\mathfrak{M}, \preceq)$ : partially ordered monoid (or set) of monomials
- $C[[\mathfrak{M}]]$ : ring of  $f = \sum_{\mathfrak{m} \in \mathfrak{M}} f_{\mathfrak{m}} \mathfrak{m}$  with **well-based support**.

$\mathfrak{m}_1 \prec \mathfrak{m}_2 \prec \dots$  with  $\mathfrak{m}_1, \mathfrak{m}_2, \dots \in \text{supp } f$  is impossible  
 $\text{supp } f$  contains no infinite antichains

- $C[\![\mathfrak{M}]\!]$ : ring of  $f = \sum_{\mathfrak{m} \in \mathfrak{M}} f_{\mathfrak{m}} \mathfrak{m}$  with **grid-based support**.

$$\text{supp } f \subseteq \mathfrak{m}_1^{\mathbb{N}} \cdots \mathfrak{m}_m^{\mathbb{N}} \{ \mathfrak{n}_1, \dots, \mathfrak{n}_n \}, \quad \mathfrak{m}_1, \dots, \mathfrak{m}_m \prec 1$$



# Examples



- $C[[z]] = C[[z^{\mathbb{N}}]] = C[[z^{\mathbb{N}}]]$
- $C((z)) = C[[z^{\mathbb{Z}}]] = C[[z^{\mathbb{Z}}]]$
- $C[[z_1, z_2]] = C[[z^{\mathbb{N}^2}]]$
- $C[[z_1]][[z_2]] = C[[z^{\mathbb{N} \times \mathbb{N}}]]$
- $C[[z^{\mathbb{Q}}]]$  Puiseux series.  $C[[z^{\mathbb{Q}}]] \supsetneq C[[z^{\mathbb{Q}}]]$



## Strong summability

We say that  $(f_i)_{i \in I} \in C[[\mathfrak{M}]]^I$  is (strongly) summable if

1.  $\bigcup_{i \in I} \text{supp } f_i$  is grid-based (or well-based if  $(f_i)_{i \in I} \in C[[\mathfrak{M}]]^I$ )
2.  $\{i \in I : \mathfrak{m} \in \text{supp } f_i\}$  is finite for each  $\mathfrak{m} \in \mathfrak{M}$

Then  $g = \sum f \in C[[\mathfrak{M}]]$  with  $g_{\mathfrak{m}} = \sum_{i \in I} f_{i, \mathfrak{m}}$  is well-defined.

## Properties

- $\sum (f_{\sigma(i)})_{i \in I} = \sum (f_i)_{i \in I}$
- $\sum F \amalg G = \sum F + \sum G$
- For  $F = \amalg_{j \in J} G_j$ , we have  $\sum_{j \in J} \sum G_j = \sum F$
- More properties



# Multiplication



## Definition

$$\left( \sum_{m \in \text{supp } f} f_m m \right) \left( \sum_{n \in \text{supp } g} g_n n \right) = \sum_{(m,n) \in \text{supp } f \times \text{supp } g} f_m g_n m n$$

- $\text{supp } f, \text{supp } g$  well/grid-based  $\Rightarrow \text{supp } f \times \text{supp } g$  well/grid-based
- $(m, n) \mapsto m n$  is increasing

## Associativity

$$f g h = \sum_{(m,n,v) \in \text{supp } f \times \text{supp } g \times \text{supp } h} f_m g_n h_v m n v$$



# Inversion



$$f = c_f \mathfrak{d}_f (1 - \varepsilon), \quad \varepsilon \prec 1$$
$$f^{-1} = c_f^{-1} \mathfrak{d}_f^{-1} \frac{1}{1 - \varepsilon}$$

$$\frac{1}{1 - \varepsilon} = \sum_{(\mathfrak{m}_1, \dots, \mathfrak{m}_l) \in (\text{supp } \varepsilon)^*} \varepsilon_{\mathfrak{m}_1} \cdots \varepsilon_{\mathfrak{m}_l} \mathfrak{m}_1 \cdots \mathfrak{m}_l$$

- Set of words  $(\text{supp } \varepsilon)^*$  carries natural partial ordering
- Higman:  $\text{supp } \varepsilon$  well-based  $\Rightarrow (\text{supp } \varepsilon)^*$  well/grid-based

More generally: for  $f \in C[[t]]$  and  $\varepsilon \prec 1$ , we may define  $f(\varepsilon)$





# Strong linearity



## Definition

Linear  $\varphi: C[[\mathfrak{M}]] \rightarrow C[[\mathfrak{N}]]$  is strongly linear if, for all summable  $(f_i)_{i \in I}$ ,  $(\varphi(f_i))_{i \in I}$  is summable and

$$\varphi\left(\sum f_i\right) = \sum \varphi(f_i)$$

## Extension by strong linearity

If  $\check{\varphi}: \mathfrak{M} \rightarrow C[[\mathfrak{N}]]$  sends grid-based subsets to summable families, then  $\check{\varphi}$  admits a unique strongly linear extension

## Applications

- For  $\varepsilon \prec 1$  and  $f, g \in C[[t]]$  with  $g_0 = 0$ ,  $f(g(\varepsilon)) = (f \circ g)(\varepsilon)$
- $(1 - \varepsilon) \frac{1}{1 - \varepsilon} = 1$



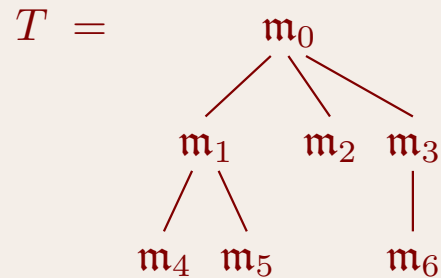
# Henselian equations



$$f = P_0 + P_1 f + P_2 f^2 + \dots, \quad f \prec 1$$

$$P_i \prec 1$$

$$f = \sum_{T \in \text{Tree}(\text{supp } f)} c_T \mathfrak{m}_T$$



$$c_T = P_{3,m_0} P_{2,m_1} P_{0,m_2} P_{1,m_3} P_{0,m_4} P_{0,m_5} P_{0,m_6}$$

$$\mathfrak{m}_T = \mathfrak{m}_0 \mathfrak{m}_1 \mathfrak{m}_2 \mathfrak{m}_3 \mathfrak{m}_4 \mathfrak{m}_5 \mathfrak{m}_6$$

Kruskal:  $\text{Tree}(\text{supp } f)$  carries well-based partial ordering



# Abstract fields of transseries



Totally ordered field  $\mathbb{T} = \mathbb{R} \llbracket \mathfrak{T} \rrbracket$  with a logarithm such that

**T1.**  $\text{dom log} = \mathbb{T}^>$ .

**T2.**  $\log m \in \mathbb{T}_{<}$ , for all  $m \in \mathfrak{T}$ , i.e.  $\forall n \in \text{supp}(\log m), n > 1$ .

**T3.**  $\log(1 + \varepsilon) = \varepsilon - \frac{1}{2} \varepsilon^2 + \frac{1}{3} \varepsilon^3 + \dots$ , for all  $\varepsilon \in \mathbb{T}_{<}$ .

**Example.**  $\mathbb{L} = \mathbb{R} \llbracket \mathfrak{L} \rrbracket = \mathbb{R} \llbracket x^{\mathbb{R}} (\log x)^{\mathbb{R}} (\log_2 x)^{\mathbb{R}} \dots \rrbracket$  with

$$\log(x^{\alpha_0} \dots (\log_k x)^{\alpha_k}) = \alpha_0 \log x + \dots + \alpha_k \log_{k+1} x$$

$$\log(f) = \log(c_f \mathfrak{d}_f (1 + \delta_f)) = \log \mathfrak{d}_f + \log c_f + \log(1 + \delta_f)$$



# The field of grid-based transseries in $x$



## Exponential extensions

$\mathbb{T} = \mathbb{R}[[\mathfrak{T}]]$  field of transseries  $\implies \mathbb{T}_{\text{exp}} = \mathbb{R}[[\mathfrak{T}_{\text{exp}}]] \supseteq \mathbb{T}$  also

$$\mathfrak{T}_{\text{exp}} = \exp(\mathbb{R}[[\mathfrak{T}]]_{>})$$

**Example.**  $e^{x^2 + \frac{x^2}{\log x} + \frac{x^2}{\log^2 x} + \dots + x + \log \log x} \in \mathfrak{L}_{\text{exp}}$

## Closure

Increasing limits of fields of grid-based transseries are fields of transseries

$$\mathbb{T} = \mathbb{L} \cup \mathbb{L}_{\text{exp}} \cup \mathbb{L}_{\text{exp,exp}} \cup \dots$$



# The field of grid-based transseries in $x$



## Exponential extensions

$\mathbb{T} = \mathbb{R}[[\mathfrak{T}]]$  field of transseries  $\implies \mathbb{T}_{\text{exp}} = \mathbb{R}[[\mathfrak{T}_{\text{exp}}]] \supseteq \mathbb{T}$  also

$$\mathfrak{T}_{\text{exp}} = \exp(\mathbb{R}[[\mathfrak{T}]]_{>})$$

## Closure fails in well-based case

$$\begin{aligned} f_1 &= x^2 \\ f_{\alpha+1} &= f_\alpha - e^{f_\alpha \circ \log x} \\ f_\lambda &= \text{stat lim}_{\alpha < \lambda} f_\alpha \end{aligned}$$



# The field of grid-based transseries in $x$



## Exponential extensions

$\mathbb{T} = \mathbb{R}[[\mathcal{T}]]$  field of transseries  $\implies \mathbb{T}_{\text{exp}} = \mathbb{R}[[\mathcal{T}_{\text{exp}}]] \supseteq \mathbb{T}$  also

$$\mathcal{T}_{\text{exp}} = \exp(\mathbb{R}[[\mathcal{T}]]_{>})$$

## Closure fails in well-based case

$$\begin{aligned}
f_1 &= x^2 \\
f_2 &= x^2 - e^{\log^2 x} \\
&\vdots \\
f_\omega &= x^2 - e^{\log^2 x} - e^{\log^2 x - e^{\log \log^2 x}} - \dots \\
f_{\omega+1} &= x^2 - e^{\log^2 x} - e^{\log^2 x - e^{\log \log^2 x}} - \dots - e^{\log^2 x - e^{\log \log^2 x} - \dots} \\
&\vdots
\end{aligned}$$



There exists a unique strong exp-log differentiation on  $\mathbb{T}$  with  $x' = 1$

**AD1.**  $f \prec g \Rightarrow f' \prec g'$ , for all  $f, g \in \mathbb{T}$  with  $g \neq 1$ .

**AD2.**  $f \succ 1 \Rightarrow (f > 0 \Rightarrow f' > 0)$ , for all  $f \in \mathbb{T}$ .

## Logarithmic transseries

- $\partial: \mathfrak{m} = x^{i_0} \cdots (\log_l x)^{i_l} \mapsto \left( \frac{i_0}{x} + \cdots + \frac{i_l}{x \cdots \log_l x} \right) \mathfrak{m}$  is grid-based
- Hence  $\partial$  admits a unique strongly linear extension on  $\mathbb{L}$

## Exponential extension of strongly linear $\partial$ on $\mathbb{T} = C[[\mathfrak{T}]]$

- $\partial: \mathfrak{m} \in \mathfrak{T}_{\text{exp}} = e^\varphi \mapsto \varphi' \mathfrak{m}$  is grid-based
- Hence  $\partial$  admits a unique strongly linear extension on  $\mathfrak{T}_{\text{exp}}$



## Composition and inversion

Given  $g \in \mathbb{T}^{>, \succ}$ , there exists a unique strong exp-log postcomposition  $\delta = \circ_g$  with  $g$  such that  $\delta x = g$  and

**AΔ1.**  $f \prec 1 \Rightarrow \delta(f) \prec 1$ , for all  $f \in \mathbb{T}$ .

**AΔ2.**  $f \geq 0 \Rightarrow \delta(f) \geq 0$ , for all  $f \in \mathbb{T}$ .

## Taylor rule

$f, \varepsilon \in \mathbb{T}$  with  $\delta \prec x$  and  $\mathfrak{m}^\dagger \delta \prec 1$  for all  $\mathfrak{m} \in \text{supp } f$ . Then

$$f \circ (x + \delta) = f + f' \delta + \frac{1}{2} f'' \delta^2 + \dots$$

## Inversion

There exists a unique  $g^{\text{inv}} \in \mathbb{T}^{>, \succ}$  with  $\delta(g^{\text{inv}}) = x$

Possible to compute  $g^{\text{inv}}$  using “Translagrange formula”





Well-based operator  $\Phi: C[[\mathfrak{M}]] \rightarrow C[[\mathfrak{M}]]$

$$\begin{aligned}\Phi &= \Phi_0 + \Phi_1 + \Phi_2 + \dots \text{ (strongly)} \\ \Phi_i(f) &= \check{\Phi}_i(f, \dots, f) \\ \check{\Phi}_i &\text{ strongly } i\text{-linear}\end{aligned}$$

## Fixed point theorem

If  $\Phi$  is *strictly extensive*, then

$$f = \Phi(f)$$

admits a unique solution in  $C[[\mathfrak{M}]]$

- Requires additional support condition in grid-based case
- Generalizes to equations  $f = \Phi(f, g)$ , obtaining  $f = \Psi(g)$



# Examples



## Functional equations

$$f = \frac{1}{x} + f(x^2) + f(e^{\log^2 x} + 1)$$

## Integration

$$\partial m = \Delta m + Rm,$$

where  $\Delta m = c_m' \partial m'$ .  $\Delta$  is strictly  $\prec$ -increasing on  $\mathfrak{T} \setminus \{1\}$ , whence  $\Delta$  and  $\Delta^{-1}$  extend by strong linearity

$$\int = \Delta^{-1} - \Delta^{-1} R \Delta^{-1} + \Delta^{-1} R \Delta^{-1} R \Delta^{-1} + \dots$$



# Asymptotic algebraic equations



Algebra

Asymptotic algebra



# Asymptotic algebraic equations



Algebra

$$P(f) = 0$$

Asymptotic algebra



# Asymptotic algebraic equations



Algebra

$$P(f) = 0$$

Asymptotic algebra

$$P(f) = 0, \quad (f \prec \mathfrak{v})$$



# Asymptotic algebraic equations



Algebra

$$P(f) = 0$$

$\deg P$

Asymptotic algebra

$$P(f) = 0, \quad (f \prec \mathfrak{v})$$



# Asymptotic algebraic equations



Algebra

$$P(f) = 0$$

$\deg P$

Asymptotic algebra

$$P(f) = 0, \quad (f \prec \mathfrak{v})$$

$\deg_{\prec \mathfrak{v}} P$



# Asymptotic algebraic equations



Algebra

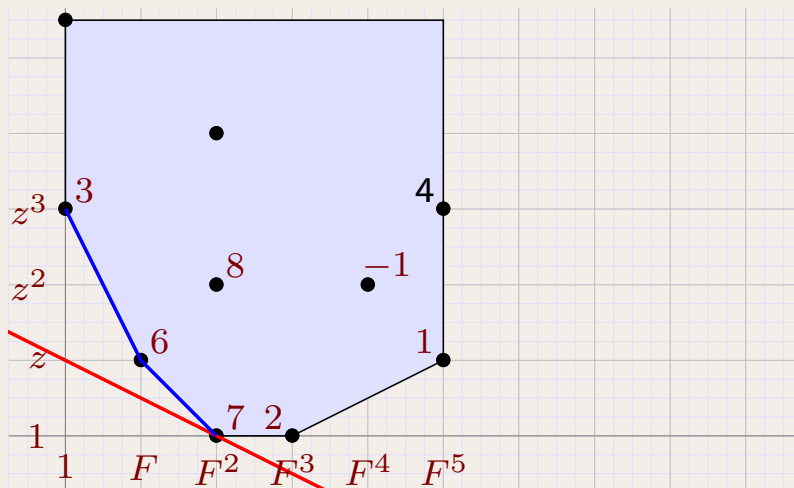
$$P(f) = 0$$

$\deg P$

Asymptotic algebra

$$P(f) = 0, \quad (f \prec \mathfrak{v})$$

$\deg_{\prec \mathfrak{v}} P$



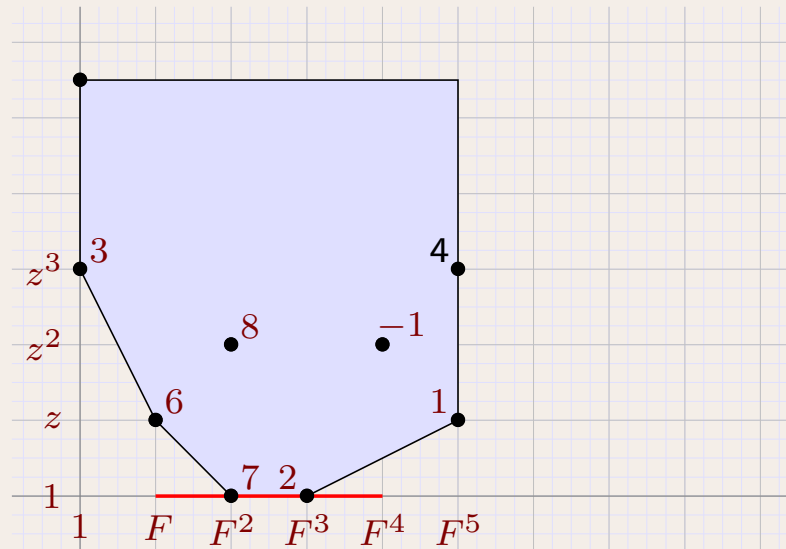




# Newton polynomials



- $P \in C[\mathcal{M}][F] \subseteq C[F][\mathcal{M}]$
- $N_P = c_P \in C[F]$



$$N_P = 2F^3 + 7F^2$$



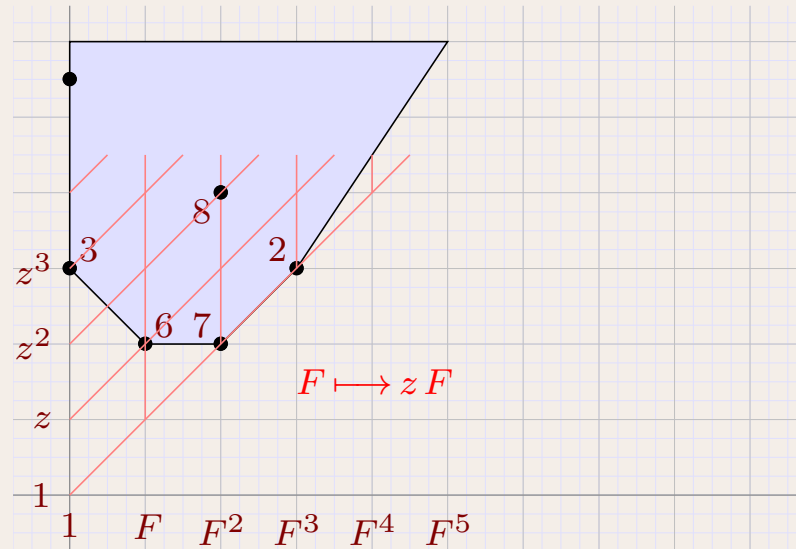
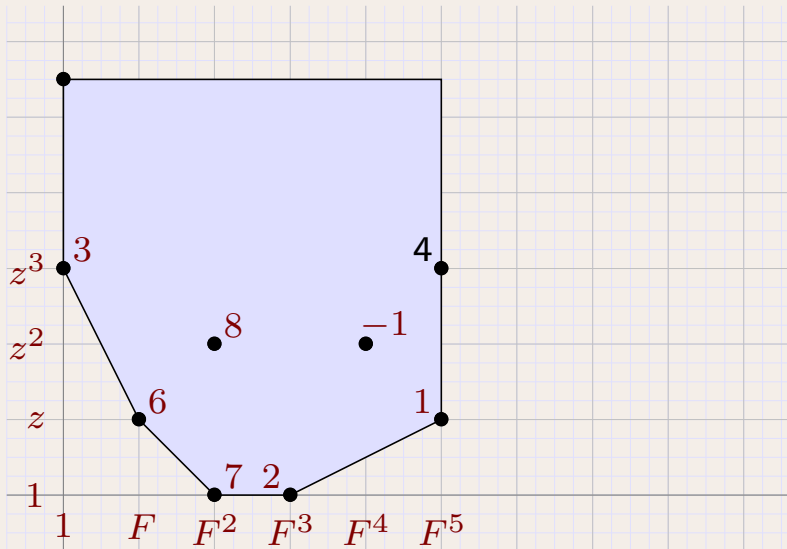
# Starting terms



- $w \prec v$  is a “starting monomial”  $\iff N_{P \times w} \notin CF^{\mathbb{N}}$
- $cw$  is a “starting term” ( $c \neq 0$ )  $\iff N_{P \times w}(c) = 0$

$$P_{\times \varphi}(f) = P(\varphi f)$$

$$P_{+\varphi}(f) = P(\varphi + f)$$





# Newton degree



$$\deg_{\prec v} P = \deg N_{P \times v}$$

$$\deg_{\prec v} P = \text{val } N_{P \times v}$$

$$\deg_{\prec w} P \leq \deg_{\prec v} P, \quad w \prec v$$

$$\deg_{\prec v} P_{+\varphi} = \deg_{\prec v} P, \quad \varphi \prec v$$

$$\deg_{\prec v} P_{\times w} = \deg_{\prec v w} P$$

$$\deg_{\prec v} (PQ) = \deg_{\prec v} P + \deg_{\prec v} Q$$

$$\deg_{\prec \varphi} P_{+\varphi} = \mu(c_\varphi; N_{P \times v_\varphi})$$

$$\mu_{\prec v}(f; P) = \deg_{\prec v} P_{+f}$$



# Newton polygon method



1.  $\deg_{\prec \mathfrak{v}} P = d > 0$   
( $P = A_{+g}$  and  $g$  root modulo  $\prec \mathfrak{v}$  of  $A$ )
2. If  $d = 1$  then unique solution
3. Determine starting monomial  $\mathfrak{w} \prec \mathfrak{v}$
4. Solve  $N_{P \times \mathfrak{w}}(c) = 0$  and set  $\varphi := c \mathfrak{w}$
5. Refine  $f = \varphi + \tilde{f}$ ,  $\tilde{f} \prec \mathfrak{w} \rightarrow 0 < \deg_{\prec \mathfrak{w}} \tilde{P} \leq d$  with  $\tilde{P} = P_{+\varphi}$   
( $\tilde{P} = A_{+g+\varphi}$  and  $g + \varphi$  root modulo  $\prec \mathfrak{w}$  of  $A$ )
6. Return to step 1



# Newton polygon method



1.  $\deg_{\prec \mathfrak{v}} P = d > 0$

( $P = A_{+g}$  and  $g$  root modulo  $\prec \mathfrak{v}$  of  $A$ )

2. If  $d = 1$  then unique solution

3. Determine starting monomial  $\mathfrak{w} \prec \mathfrak{v}$

4. Solve  $N_{P \times \mathfrak{w}}(c) = 0$  and set  $\varphi := c \mathfrak{w}$

$$\left( f - \frac{1}{1-z} \right)^2 = z^{10000}$$

5. Refine  $f = \varphi + \tilde{f}$ ,  $\tilde{f} \prec \mathfrak{w} \rightarrow 0 < \deg_{\prec \mathfrak{w}} \tilde{P} \leq d$  with  $\tilde{P} = P_{+\varphi}$

( $\tilde{P} = A_{+g+\varphi}$  and  $g + \varphi$  root modulo  $\prec \mathfrak{w}$  of  $A$ )

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# Newton polygon method



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3. Determine starting monomial  $\mathfrak{w} \prec \mathfrak{v}$

4. Solve  $N_{P \times \mathfrak{w}}(c) = 0$  and set  $\varphi := c \mathfrak{w}$

**If  $\mu_{N_{P \times \mathfrak{w}}}(c) = d$ , then  $\varphi :=$  unique solution to  $\frac{\partial^{d-1} P}{\partial F^{d-1}}(\varphi) = 0$ ,  $\varphi \prec \mathfrak{v}$**

5. Refine  $f = \varphi + \tilde{f}$ ,  $\tilde{f} \prec \mathfrak{w} \longrightarrow 0 < \deg_{\prec \mathfrak{w}} \tilde{P} \leq d$  with  $\tilde{P} = P_{+\varphi}$

( $\tilde{P} = A_{+g+\varphi}$  and  $g + \varphi$  root modulo  $\prec \mathfrak{w}$  of  $A$ )

6. Return to step 1

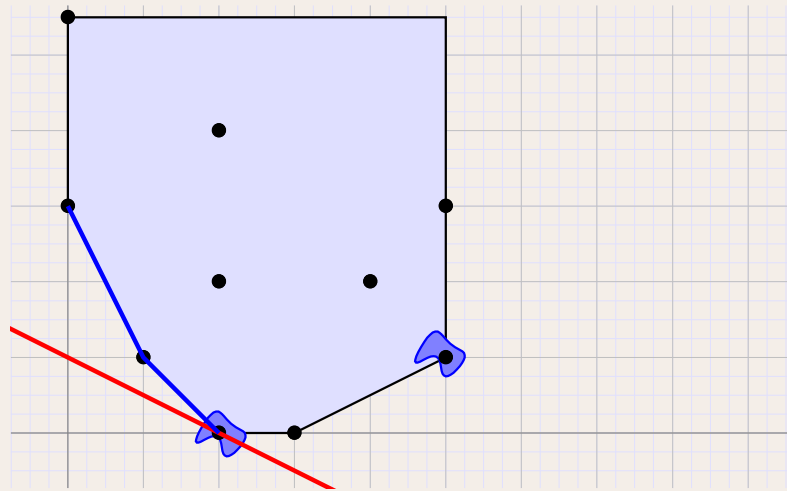


# Differential Newton polygon method



$$P(f) = p(f, f', \dots, f^{(r)}) = 0, \quad f \prec \mathfrak{v}$$

Starting monomials cannot directly be read off from “Newton polygon”



$$P = P_0 + \dots + P_d$$



# Upward shifting



$P\uparrow$  unique differential polynomial with

$$(P\uparrow)(f \circ e^x) = P(f) \circ e^x$$

For instance:

$$\begin{aligned} F'\uparrow &= \frac{F'}{e^x} \\ F''\uparrow &= \frac{F'' - F'}{e^{2x}} \\ F'''\uparrow &= \frac{F''' - 3F'' + 2F'}{e^{3x}} \\ &\vdots \end{aligned}$$





# Differential Newton polynomial



**Theorem.** *There exists a unique  $N_P \in \mathbb{R}\{F\}$ , such that*

$$c_{P \uparrow l} = N_P$$

*for all sufficiently large  $l$  and*

$$N_P \in \mathbb{R}[F] (F')^{\mathbb{N}}.$$

**Definition.**  $\mathfrak{m} \prec \mathfrak{v}$  *is a starting monomial*  $\iff N_{P \times \mathfrak{m}} \notin \mathbb{R} F^{\mathbb{N}}$



## Example



$$\begin{aligned}P &= (F')^2 - FF'' \\P\uparrow &= \frac{(F')^2 - FF'' + FF'}{e^{2x}} \\P\uparrow\uparrow &= \frac{FF'}{e^x e^{2e^x}} + \frac{(F')^2 - FF'' + FF'}{e^{2x} e^{2e^x}} \\&\vdots \\N_P &= FF'\end{aligned}$$

Consequence:

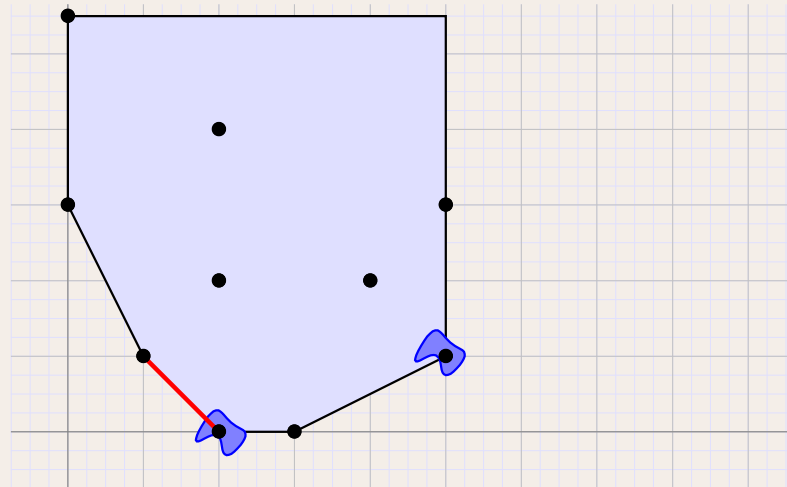
$$1 \prec L \prec \log_n x \implies P(L) \sim \frac{LL'}{x}$$



# Starting monomials



**Lemma.** Given  $i < j$  with  $P_i \neq 0$ ,  $P_j \neq 0$ , there exists a unique  $(i, j)$ -equalizer  $\mathfrak{e} \in \mathfrak{T}$  such that  $N_{(P_i+P_j) \times \mathfrak{e}}$  is not homogeneous.





# Starting monomials

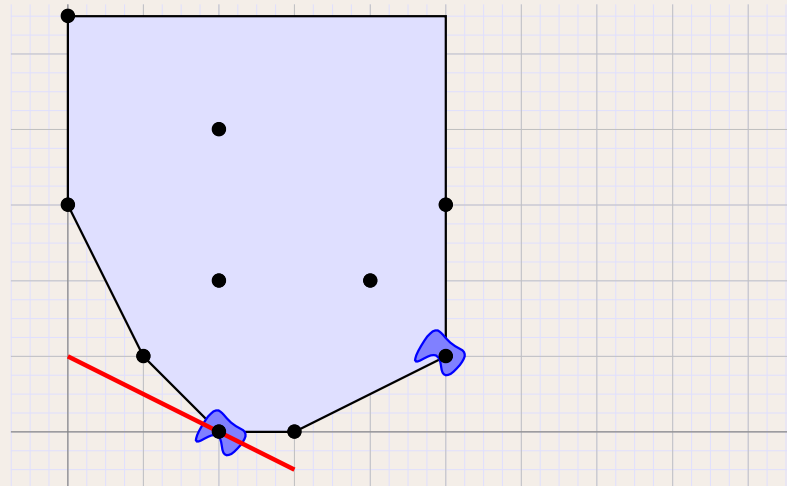


**Lemma.** Given  $i$  with  $P_i \neq 0$ , we have

$\mathfrak{m}$  is a starting monomial for  $P_i(f) = 0$



$\mathfrak{m}^\dagger = \frac{\mathfrak{m}'}{\mathfrak{m}}$  is a solution to  $R_{P_i}(g) = 0$  modulo  $\frac{1}{x \log x \log_2 x \cdots}$





# Solving asymptotic differential equations



**Lemma.**  $\deg_{\prec v} P = 1 \implies P(f) = 0, f \prec v$  admits at least one solution.

**Warning.** Problem with almost multiple solutions

$$f^2 - 2 f' + \frac{1}{x^2} + \dots + \frac{1}{(x \log x \dots \log_l x)^2} = 0, \quad (f \prec 1)$$

$$f^2 - 2 e^{-x} f' + \frac{1}{e^{2x}} + \dots + \frac{1}{(e^x x \dots \log_{l-1} x)^2} = 0, \quad (f \prec 1)$$

$$f^2 - 2 f' - 2 f + 1 + \frac{1}{x^2} + \dots + \frac{1}{(x \log x \dots \log_{l-1} x)^2} = 0, \quad (f \prec 1)$$

$$f^2 - 2 f' + \frac{1}{x^2} + \dots + \frac{1}{(x \log x \dots \log_{l-1} x)^2} = 0, \quad (f \prec 1)$$

**Lemma.** “Unravelling process” is finite for grid-based transseries.



# Results



**Theorem.** (1997) *There exists a theoretical algorithm to find all solutions to an asymptotic algebraic differential equation.*

**Theorem.** (1997) *Let  $P$  be purely exponential of degree  $d$  and order  $r$ . There exists a constant  $C_{r,d}$  such that any solution to  $P(f) = 0$  involves at most  $C_{r,d}$  levels of iterated logarithms.*

**Theorem.** (1997) *Any general transseries solution to an algebraic differential equation with grid-based coefficients is again grid-based. Generalization of Grigoriev and Singer (1991).*

**Corollary.**  $\zeta(x)$  and  $f(x) = \frac{1}{x} + \frac{1}{e^{\log^2 x}} + \frac{1}{e^{\log^4 x}} + \dots$  are differentially transcendental over  $\mathbb{R}$ .