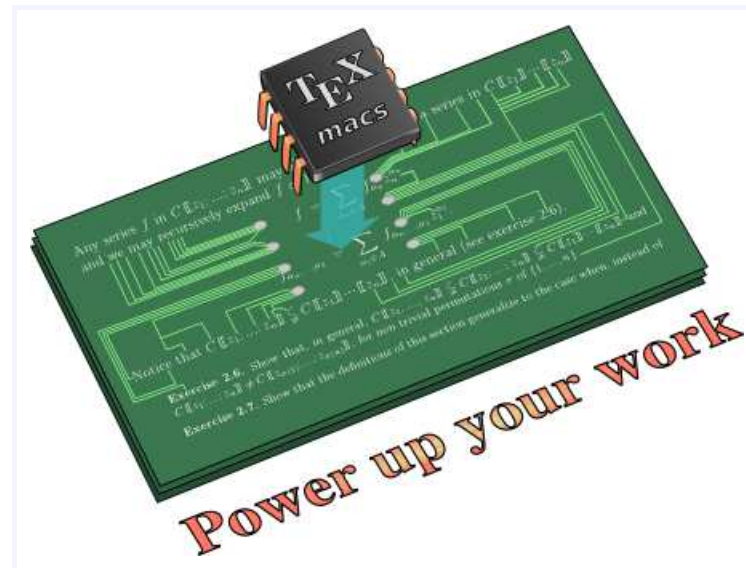


Sparse polynomial interpolation in practice

Joris van der Hoeven, Grégoire Lecerf

CNRS, École polytechnique



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Motivation (implementation in upcoming CAAS package, matrices)

```
Caas] M(n) == [ x[i,j] | i in 1 to n || j in 1 to n ];
```

```
Caas] M(2)
```

$$\begin{bmatrix} x_{1,1} & x_{2,1} \\ x_{1,2} & x_{2,2} \end{bmatrix}$$

Caas] invert M(2)

$$\begin{bmatrix} \frac{x_{2,1} x_{1,2}}{\left(x_{2,2} - \frac{x_{2,1} x_{1,2}}{x_{1,1}}\right) x_{1,1}^2} + \frac{1}{x_{1,1}} & -\left(\frac{x_{2,1}}{\left(x_{2,2} - \frac{x_{2,1} x_{1,2}}{x_{1,1}}\right) x_{1,1}}\right) \\ -\left(\frac{x_{1,2}}{\left(x_{2,2} - \frac{x_{2,1} x_{1,2}}{x_{1,1}}\right) x_{1,1}}\right) & \frac{1}{x_{2,2} - \frac{x_{2,1} x_{1,2}}{x_{1,1}}} \end{bmatrix}$$

Caas] simplify (M(12) * invert M(12))

$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Caas] simplify transpose invert transpose invert M(8)

$$\begin{bmatrix} x_{1,1} & x_{2,1} & x_{3,1} & x_{4,1} & x_{5,1} & x_{6,1} & x_{7,1} & x_{8,1} \\ x_{1,2} & x_{2,2} & x_{3,2} & x_{4,2} & x_{5,2} & x_{6,2} & x_{7,2} & x_{8,2} \\ x_{1,3} & x_{2,3} & x_{3,3} & x_{4,3} & x_{5,3} & x_{6,3} & x_{7,3} & x_{8,3} \\ x_{1,4} & x_{2,4} & x_{3,4} & x_{4,4} & x_{5,4} & x_{6,4} & x_{7,4} & x_{8,4} \\ x_{1,5} & x_{2,5} & x_{3,5} & x_{4,5} & x_{5,5} & x_{6,5} & x_{7,5} & x_{8,5} \\ x_{1,6} & x_{2,6} & x_{3,6} & x_{4,6} & x_{5,6} & x_{6,6} & x_{7,6} & x_{8,6} \\ x_{1,7} & x_{2,7} & x_{3,7} & x_{4,7} & x_{5,7} & x_{6,7} & x_{7,7} & x_{8,7} \\ x_{1,8} & x_{2,8} & x_{3,8} & x_{4,8} & x_{5,8} & x_{6,8} & x_{7,8} & x_{8,8} \end{bmatrix}$$

Caas] simplify invert M(2)

$$\begin{bmatrix} \frac{x_{2,2}}{x_{2,2}x_{1,1} - x_{2,1}x_{1,2}} & -\left(\frac{x_{2,1}}{x_{2,2}x_{1,1} - x_{2,1}x_{1,2}}\right) \\ -\left(\frac{x_{1,2}}{x_{2,2}x_{1,1} - x_{2,1}x_{1,2}}\right) & \frac{x_{1,1}}{x_{2,2}x_{1,1} - x_{2,1}x_{1,2}} \end{bmatrix}$$

Caas] simplify invert M(3)

$$\begin{bmatrix} \frac{x_{3,3}x_{2,2} - x_{3,2}x_{2,3}}{(x_{3,3}x_{2,2} - x_{3,2}x_{2,3})x_{1,1} + (x_{3,2}x_{2,1} - x_{3,1}x_{2,2})x_{1,3} + (x_{3,1}x_{2,3} - x_{3,3}x_{2,1})x_{1,2}} & \frac{x_{3,2}x_{1,3} - x_{3,3}x_{1,2}}{(x_{3,3}x_{2,2} - x_{3,2}x_{2,3})x_{1,1} + (x_{3,2}x_{2,1} - x_{3,1}x_{2,2})x_{1,3} + (x_{3,1}x_{2,3} - x_{3,3}x_{2,1})x_{1,2}} \\ \frac{x_{2,3}x_{1,2} - x_{2,2}x_{1,3}}{(x_{3,3}x_{2,2} - x_{3,2}x_{2,3})x_{1,1} + (x_{3,2}x_{2,1} - x_{3,1}x_{2,2})x_{1,3} + (x_{3,1}x_{2,3} - x_{3,3}x_{2,1})x_{1,2}} & \frac{x_{3,3}x_{2,2} - x_{3,2}x_{2,3}}{(x_{3,3}x_{2,2} - x_{3,2}x_{2,3})x_{1,1} + (x_{3,2}x_{2,1} - x_{3,1}x_{2,2})x_{1,3} + (x_{3,1}x_{2,3} - x_{3,3}x_{2,1})x_{1,2}} \end{bmatrix}$$

Caas]

Caas] e1 == exp series (0, x + y)

$$1 + (x + y) z + \frac{1}{2} (x + y)^2 z^2 + \frac{1}{6} (x + y)^3 z^3 + \frac{1}{24} (x + y)^4 z^4 + \frac{1}{120} (x + y)^5 z^5 + \frac{1}{720} (x + y)^6 z^6 + \frac{1}{5040} (x + y)^7 z^7 + \frac{1}{40320} (x + y)^8 z^8 + \frac{1}{362880} (x + y)^9 z^9 + O(z^{10})$$

Caas] e2 == exp series (0, x)

$$1 + x z + \frac{1}{2} x^2 z^2 + \frac{1}{6} x^3 z^3 + \frac{1}{24} x^4 z^4 + \frac{1}{120} x^5 z^5 + \frac{1}{720} x^6 z^6 + \frac{1}{5040} x^7 z^7 + \frac{1}{40320} x^8 z^8 + \frac{1}{362880} x^9 z^9 + O(z^{10})$$

Caas] e1 / e2

Caas] q == simplify (e1 / e2)

$$1 + y z + \frac{1}{2} y^2 z^2 + \frac{1}{6} y^3 z^3 + \frac{1}{24} y^4 z^4 + \frac{1}{120} y^5 z^5 + \frac{1}{720} y^6 z^6 + \frac{1}{5040} y^7 z^7 + \frac{1}{40320} y^8 z^8 + \frac{1}{362880} y^9 z^9 + O(z^{10})$$

Caas] q [15]

$$\frac{1}{1307674368000} y^{15}$$

\mathbb{K} field (usually $\mathbb{K} = \mathbb{Q}$ or $\mathbb{K} = \mathbb{F}_p$)

Input: an expression P^{expr} in $x_1, \dots, x_n, +, -, \times, \div$ and constants in \mathbb{K}

Output: a sparse polynomial $P = \sum_{i=1}^t c_i x^{e_i}$ representing P^{expr} resp. P^{fun}

Optionally: a bound on t , or on $d = \deg P$, or on the $d_i = \deg_{x_i} P$

Note: in absence of a bound on t , we are satisfied with probabilistic algorithms

For instance, picking $x := \alpha$ random, we have $P^{\text{fun}}(\alpha) = 0 \Leftrightarrow P = 0$ with high probability

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Output: $\left\{ \begin{array}{l} \text{a sparse polynomial } P = \sum_{i=1}^t c_i x^{e_i} \text{ representing } P^{\text{expr}} \text{ resp. } P^{\text{fun}}, \\ \text{where } c_i \in \mathbb{K} \setminus \{0\}, e_i \in \mathbb{N}^n \text{ pairwise distinct; or } \perp \end{array} \right.$

Optionally: a bound on t , or on $d = \deg P$, or on the $d_i = \deg_{x_i} P$

Note: in absence of a bound on t , we are satisfied with probabilistic algorithms

For instance, picking $x := \alpha$ random, we have $P^{\text{fun}}(\alpha) = 0 \Leftrightarrow P = 0$ with high probability

Caas] $M == [a, b; c, d]$

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix}$$

Caas] $M * \text{invert } M$

$$\begin{bmatrix} \left(\frac{bc}{\left(d - \frac{bc}{a}\right)a^2} + \frac{1}{a} \right) a - \frac{bc}{\left(d - \frac{bc}{a}\right)a} & 0 \\ \left(\frac{bc}{\left(d - \frac{bc}{a}\right)a^2} + \frac{1}{a} \right) c - \frac{cd}{\left(d - \frac{bc}{a}\right)a} & \frac{d}{d - \frac{bc}{a}} - \frac{bc}{\left(d - \frac{bc}{a}\right)a} \end{bmatrix}$$

Caas]

Caas] M == [force_simplify_zero a, b; c, d]

$$\begin{bmatrix} a & b \\ c & d \end{bmatrix}$$

Caas] M * invert M

$$\begin{bmatrix} \left(\frac{bc}{\left(d - \frac{bc}{a}\right)a^2} + \frac{1}{a} \right) a - \frac{bc}{\left(d - \frac{bc}{a}\right)a} & 0 \\ 0 & \frac{d}{d - \frac{bc}{a}} - \frac{bc}{\left(d - \frac{bc}{a}\right)a} \end{bmatrix}$$

Caas]

Assume (for known t)

$$P = \sum_{i=1}^t c_i x^{e_i}.$$

For all $\alpha \in \mathbb{K}^n$ we have

$$f(z) := \sum_{i \geq 0} P(\alpha_1^i, \dots, \alpha_n^i) z^i = \sum_{1 \leq i \leq t} \frac{c_i}{1 - \alpha^{e_i} z}.$$

Choose α such that the α^{e_i} are pairwise distinct

Compute $c_1, \alpha^{e_1}, \dots, c_t, \alpha^{e_t} \in \mathbb{K}$ from f_0, \dots, f_{2t-1}

- Find $A, B \in \mathbb{K}[z]$ of degrees $< t$ with $f(z) = A(z)/B(z)$ using Padé-Hermite
- Find the roots $\alpha^{e_1}, \dots, \alpha^{e_t}$ of B
- Solve linear Vandermonde system in order to get the c_i

Recover e_1, \dots, e_t from $\alpha^{e_1}, \dots, \alpha^{e_t}$

Without bound for t , try for $t = 1, 2, 4, 8, 16, \dots$

Prime numbers. [Ben-Or & Tiwari 1988], \mathbb{K} characteristic zero

Take $\alpha_i = \pi_i$ (the i -th prime number)

Roots of integer polynomials can be computed fast

$\alpha^{e_i} \rightsquigarrow e_i$ via factoring

Modular arithmetic. [Kaltofen & Lakshman & Wiley 1990], DAGWOOD

Pick a prime p such that $\max_i \alpha^{e_i} < p$ (e.g. $\pi_n^d < p$)

Evaluate P using arithmetic modulo p only

Find roots of polynomials in \mathbb{F}_p using Cantor-Zassenhaus or Moenck (if $p = k 2^l + 1$)

$\alpha^{e_i} \rightsquigarrow e_i$ via factoring

Kronecker approach. [Javadi & Monagan 2010], [Arnold & Roche 2014]

Let $D_i \geq \deg_{x_i} P + 1$

Let $p > D_1 \cdots D_n$ and let ω be a primitive element in \mathbb{F}_p

Take $\alpha = (1, \omega^{D_1}, \omega^{D_1 D_2}, \dots, \omega^{D_1 \cdots D_{n-1}})$

Find roots of polynomials in \mathbb{F}_p using Cantor-Zassenhaus or Moenck

$\alpha^{e_i} = \omega^{\Pi \cdot e_i} \rightsquigarrow e_i$ via discrete logarithm (fast if $p = k 2^l + 1$)

Let $Q(x_{m+1}, \dots, x_n) = P(\alpha_1, \dots, \alpha_m, x_{m+1}, \dots, x_n)$ for random α

Problem: assuming $\text{supp } Q = \{f_1, \dots, f_u\}$ known, efficiently recover $\text{supp } P$

Hope: possible using a much smaller p (fitting in one word)

Coefficient ratios. [Javadi & Monagan, slightly different context], $\pi_1^{e_{i,1}} \dots \pi_m^{e_{i,m}} < p$

$$\begin{aligned} f(x_1, \dots, x_n) &\rightsquigarrow c_i, \alpha^{e_i} \\ f(\pi_1 x_1, \dots, \pi_m x_m, x_{m+1}, \dots, x_n) &\rightsquigarrow \pi_1^{e_{i,1}} \dots \pi_m^{e_{i,m}} c_i, \alpha^{e_i} \rightsquigarrow e_{i, \leq m} := (e_{i,1}, \dots, e_{i,m}) \\ \alpha^{e_i - e_{i, \leq m}}, \text{supp } Q &\rightsquigarrow e_i \end{aligned}$$

Works also for Kronecker approach

Closest point projections. $\alpha = (1, \omega^{D_1}, \dots, \omega^{D_1 \dots D'_m \dots D_{n-1}})$, $D'_m > D_m$ random

Consider $\varphi_i \in \{0, \dots, p-2\}$ with $\alpha^{f_i} = \omega^{\varphi_i}$

The φ_i modulo $p-1$ are essentially random with minimal distance $\delta \approx p/u^2$

Assume $D_1 \dots D_m < \delta$

For any $e_i \in \text{supp } P$ and ψ_i with $\alpha^{e_i} = \omega^{\psi_i}$, $\exists! \varphi_{i'}$ with $\psi_i - \varphi_{i'} < \delta$ (modulo $p-1$)

Problem: interpolate generic $k \times k$ determinant

$$P = \begin{vmatrix} x_{1,1} & \cdots & x_{1,k} \\ \vdots & & \vdots \\ x_{k,1} & \cdots & x_{k,k} \end{vmatrix}$$

Interesting parameters, since

$$n = k^2$$

$$t = k!$$

$$d = k$$

$$d_i = 1$$

n	4	5	6	7	8	9	10	11	12
Bound $\log_2 p$ (primes)	23	33	44	55	67	79	91	104	117
Bound $\log_2 p$ (Kronecker)	16	25	36	49	64	81	100	121	144

```
Mmx] use "multimix";
```

```
Mmx] n == 4;
```

```
Mmx] M == [ polynomial_dag (1 := Integer, coordinate ('x[i,j])) |  
            i in 1 to n || j in 1 to n ]
```

$$\begin{bmatrix} x_{1,1} & x_{2,1} & x_{3,1} & x_{4,1} \\ x_{1,2} & x_{2,2} & x_{3,2} & x_{4,2} \\ x_{1,3} & x_{2,3} & x_{3,3} & x_{4,3} \\ x_{1,4} & x_{2,4} & x_{3,4} & x_{4,4} \end{bmatrix}$$

```
Mmx] D == det M
```

$$\begin{aligned} & -(x_{1,4} x_{4,1} + x_{2,4} x_{4,2} + x_{3,4} x_{4,3}) (-x_{1,3} x_{3,1} - x_{2,3} x_{3,2} - x_{3,3} (-x_{2,2} - x_{1,1}) - x_{1,2} x_{2,1} + \\ & x_{2,2} x_{1,1}) - (x_{1,4} (x_{1,1} x_{4,1} + x_{2,1} x_{4,2} + x_{3,1} x_{4,3}) + x_{2,4} (x_{1,2} x_{4,1} + x_{2,2} x_{4,2} + x_{3,2} x_{4,3}) + \\ & x_{3,4} (x_{1,3} x_{4,1} + x_{2,3} x_{4,2} + x_{3,3} x_{4,3})) (-x_{3,3} - x_{2,2} - x_{1,1}) - x_{1,4} (x_{1,1} (x_{1,1} x_{4,1} + x_{2,1} x_{4,2} + \\ & x_{3,1} x_{4,3}) + x_{2,1} (x_{1,2} x_{4,1} + x_{2,2} x_{4,2} + x_{3,2} x_{4,3}) + x_{3,1} (x_{1,3} x_{4,1} + x_{2,3} x_{4,2} + x_{3,3} x_{4,3})) - \\ & x_{2,4} (x_{1,2} (x_{1,1} x_{4,1} + x_{2,1} x_{4,2} + x_{3,1} x_{4,3}) + x_{2,2} (x_{1,2} x_{4,1} + x_{2,2} x_{4,2} + x_{3,2} x_{4,3}) + x_{3,2} (x_{1,3} x_{4,1} + \\ & x_{2,3} x_{4,2} + x_{3,3} x_{4,3})) - x_{3,4} (x_{1,3} (x_{1,1} x_{4,1} + x_{2,1} x_{4,2} + x_{3,1} x_{4,3}) + x_{2,3} (x_{1,2} x_{4,1} + x_{2,2} x_{4,2} + \\ & x_{3,2} x_{4,3}) + x_{3,3} (x_{1,3} x_{4,1} + x_{2,3} x_{4,2} + x_{3,3} x_{4,3})) - x_{4,4} (- (x_{1,3} x_{3,1} + x_{2,3} x_{3,2}) (-x_{2,2} - x_{1,1}) - \\ & x_{1,3} (x_{1,1} x_{3,1} + x_{2,1} x_{3,2}) - x_{2,3} (x_{1,2} x_{3,1} + x_{2,2} x_{3,2}) - x_{3,3} (-x_{1,2} x_{2,1} + x_{2,2} x_{1,1})) \end{aligned}$$

```
Mmx] as_mvpolynomial% (D)
```

$$\begin{aligned} & x_{1,1} x_{2,2} x_{3,3} x_{4,4} - x_{1,2} x_{2,1} x_{3,3} x_{4,4} - x_{1,1} x_{2,3} x_{3,2} x_{4,4} + x_{1,3} x_{2,1} x_{3,2} x_{4,4} + x_{1,2} x_{2,3} x_{3,1} x_{4,4} - \\ & x_{1,3} x_{2,2} x_{3,1} x_{4,4} - x_{1,1} x_{2,2} x_{3,4} x_{4,3} + x_{1,2} x_{2,1} x_{3,4} x_{4,3} + x_{1,1} x_{2,4} x_{3,2} x_{4,3} - x_{1,4} x_{2,1} x_{3,2} x_{4,3} - \\ & x_{1,2} x_{2,4} x_{3,1} x_{4,3} + x_{1,4} x_{2,2} x_{3,1} x_{4,3} + x_{1,1} x_{2,3} x_{3,4} x_{4,2} - x_{1,3} x_{2,1} x_{3,4} x_{4,2} - x_{1,1} x_{2,4} x_{3,3} x_{4,2} + \\ & x_{1,4} x_{2,1} x_{3,3} x_{4,2} + x_{1,3} x_{2,4} x_{3,1} x_{4,2} - x_{1,4} x_{2,3} x_{3,1} x_{4,2} - x_{1,2} x_{2,3} x_{3,4} x_{4,1} + x_{1,3} x_{2,2} x_{3,4} x_{4,1} + \\ & x_{1,2} x_{2,4} x_{3,3} x_{4,1} - x_{1,4} x_{2,2} x_{3,3} x_{4,1} - x_{1,3} x_{2,4} x_{3,2} x_{4,1} + x_{1,4} x_{2,3} x_{3,2} x_{4,1} \end{aligned}$$


```
template<int n> integer
det_modulo (const vector<integer>& v /* n x n matrix coefficients */,
           const integer& p);
```

```
template<int n> void
print_generic_determinant () {
    function_2<integer, const vector<integer>&, const integer&>
        f (&det_modulo<n>);
    sparse_polynomial<integer> pol =
        pr_as_sparse_polynomial (f, n*n, n /* degree bound */);
    mmout << pol << lf;
}
```

n	5	6	7	8
Kronecker substitution	110 ms	915 ms	9.8 s	162 s

Table 1. Timings for interpolating the determinant polynomial of a $n \times n$ matrix.